

# 1 Syntax

$e ::= c \in \mathbb{Z}$	$e_e ::= \cdot +_1 e$	$s \in stat ::= skip$	$s_e ::= x :=_1 \cdot$
$\quad   x \in Var$	$\quad   \cdot +_2 \cdot$	$\quad   s_1; s_2$	$\quad   \cdot ;_1 s_2$
$\quad   e_1 + e_2$	$\quad   @_1(e_2)$	$\quad   x := e$	$\quad   if_1 s_1 s_2$
$\quad   \lambda x. s$	$\quad   @_2$	$\quad   if (e > 0) s_1 s_2$	$\quad   while_1 (e > 0) s$
$\quad   e_1(e_2)$	$\quad   @_3$	$\quad   while (e > 0) s$	$\quad   while_2 (e > 0) s$
		$\quad   return e$	$\quad   return_1 \cdot$

# 2 Semantics

## 2.1 Expressions

$$\frac{\text{RED-CONST}(c)}{H_e, \ell_e, \ell_c, c \Downarrow H_e, \ell_e, c} \quad \frac{\text{RED-VAR-LOCAL}(x)}{H_e, \ell_e, \ell_c, x \Downarrow H_e, \ell_e, \ell_c[x]} \quad x \in \text{dom}(H_e[\ell_c])$$

$$\frac{\text{RED-VAR-GLOBAL}(x)}{H_e, \ell_e, \ell_c, x \Downarrow H_e, \ell_e, E[x]} \quad x \in \text{dom}(H_e[\ell_e]) \wedge x \notin \text{dom}(H_e[\ell_c])$$

$$\frac{\text{RED-VAR-UNDEF}(x)}{H_e, \ell_e, \ell_c, x \Downarrow err} \quad x \notin \text{dom}(H_e[\ell_e]) \wedge x \notin \text{dom}(H_e[\ell_c])$$

$$\frac{\text{RED-ADD}(e_1, e_2)}{H_e, \ell_e, \ell_c, e_1 \Downarrow r \quad \ell_c, r, \cdot +_1 e_2 \Downarrow r'}{H_e, \ell_e, \ell_c, e_1 + e_2 \Downarrow r'} \quad \frac{\text{RED-ADD-1}(e_2)}{H_e, \ell_e, \ell_c, e_2 \Downarrow r \quad v_1, r, \cdot +_2 \cdot \Downarrow r'}{\ell_c, (H_e, \ell_e, v_1), \cdot +_1 e_2 \Downarrow r'}$$

$$\frac{\text{RED-ADD-2}}{v_1, (H_e, \ell_e, v_2), \cdot +_2 \cdot \Downarrow H_e, \ell_e, v_1 + v_2} \quad \frac{\text{RED-LAMBDA}(x, s)}{H_e, \ell_e, \ell_c, \lambda x. s \Downarrow H_e, \ell_e, (\ell_c, \lambda x. s)}$$

$$\frac{\text{RED-APP}(e_1, e_2)}{H_e, \ell_e, \ell_c, e_1 \Downarrow r \quad \ell_c, r, @_1(e_2) \Downarrow r'}{H_e, \ell_e, \ell_c, e_1(e_2) \Downarrow r'} \quad \frac{\text{RED-APP-1}(e_2)}{H_e, \ell_e, \ell_c, e_2 \Downarrow r \quad \ell'_c, x, s, r, @_2 \Downarrow r'}{\ell_c, (H_e, \ell_e, (\ell'_c, \lambda x. s)), @_1(e_2) \Downarrow r'}$$

$$\frac{\text{RED-APP-2}(s)}{H_e[\ell'_c \leftarrow C[x \leftarrow v]], \ell_e, \ell'_c, s \Downarrow r \quad r, @_3 \Downarrow r'}{\ell_c, x, s, (H_e, \ell_e, v), @_2 \Downarrow r'} \quad \frac{\text{RED-APP-3-RET}}{ret(H_e, \ell_e, v), @_3 \Downarrow H_e, \ell_e, v}$$

$$\frac{\text{RED-APP-3-NO-RET}}{H_e, \ell_e, \ell_c, @_3 \Downarrow err}$$

## 2.2 Statements

$$\begin{array}{c}
\text{RED-SKIP} \\
\frac{}{H_e, \ell_e, \ell_c, \text{skip} \Downarrow H_e, \ell_e, \ell_c}
\end{array}
\qquad
\frac{\text{RED-SEQ}(s_1, s_2)}{H_e, \ell_e, \ell_c, s_1 \Downarrow r \quad r, \cdot;_1 s_2 \Downarrow r'}{H_e, \ell_e, \ell_c, s_1; s_2 \Downarrow r'}$$

$$\frac{\text{RED-SEQ-1}(s_2)}{H_e, \ell_e, \ell_c, s_2 \Downarrow r}{H_e, \ell_e, \ell_c, \cdot;_1 s_2 \Downarrow r}
\qquad
\frac{\text{RED-ASN}(x, e)}{H_e, \ell_e, \ell_c, e \Downarrow r \quad \ell_c, r, x :=_1 \cdot \Downarrow r'}{H_e, \ell_e, \ell_c, x := e \Downarrow r'}$$

$$\frac{\text{RED-ASN-1}(x)}{\ell'_e = \text{fresh}(H_e) \quad E = H_e[\ell_e]}{\ell_c, (H_e, \ell_e, v), x :=_1 \cdot \Downarrow H_e[\ell'_e \leftarrow E[x \leftarrow v]], \ell'_e, \ell_c} \quad x \notin \text{dom}(H_e[\ell_c])$$

$$\frac{\text{RED-ASN-1-LOCAL}(x)}{\ell'_c = \text{fresh}(H_e) \quad C = H_e[\ell_c]}{\ell_c, (H_e, \ell_e, v), x :=_1 \cdot \Downarrow H_e[\ell'_c \leftarrow C[x \leftarrow v]], \ell_e, \ell'_c} \quad x \in \text{dom}(C)$$

$$\frac{\text{RED-IF}(e, s_1, s_2)}{H_e, \ell_e, \ell_c, e \Downarrow r \quad \ell_c, r, \text{if}_1 s_1 s_2 \Downarrow r'}{H_e, \ell_e, \ell_c, \text{if}(e > 0) s_1 s_2 \Downarrow r'}
\qquad
\frac{\text{RED-IF-1-POS}(s_1, s_2)}{H_e, \ell_e, \ell_c, s_1 \Downarrow r}{\ell_c, (H_e, \ell_e, v), \text{if}_1 s_1 s_2 \Downarrow r} \quad v > 0$$

$$\frac{\text{RED-IF-1-NEG}(s_1, s_2)}{H_e, \ell_e, \ell_c, s_2 \Downarrow r}{\ell_c, (H_e, \ell_e, v), \text{if}_1 s_1 s_2 \Downarrow r} \quad v \leq 0$$

$$\frac{\text{RED-WHILE}(e, s)}{H_e, \ell_e, \ell_c, e \Downarrow r \quad \ell_c, r, \text{while}_1(e > 0) s \Downarrow r'}{H_e, \ell_e, \ell_c, \text{while}(e > 0) s \Downarrow r'}$$

$$\frac{\text{RED-WHILE-1-NEG}(e, s)}{\ell_c, (H_e, \ell_e, v), \text{while}_1(e > 0) s \Downarrow H_e, \ell_e, \ell_c} \quad v \leq 0$$

$$\frac{\text{RED-WHILE-1-POS}(e, s)}{H_e, \ell_e, \ell_c, s \Downarrow r \quad r, \text{while}_2(e > 0) s \Downarrow r'}{\ell_c, (H_e, \ell_e, v), \text{while}_1(e > 0) s \Downarrow r'} \quad v > 0$$

$$\frac{\text{RED-WHILE-2}(e, s)}{H_e, \ell_e, \ell_c, \text{while}(e > 0) s \Downarrow r}{H_e, \ell_e, \ell_c, \text{while}_2(e > 0) s \Downarrow r}
\qquad
\frac{\text{RED-RETURN}(e)}{H_e, \ell_e, \ell_c, e \Downarrow r \quad r, \text{return}_1 \cdot \Downarrow r'}{H_e, \ell_e, \ell_c, \text{return } e \Downarrow r'}$$

$$\frac{\text{RED-RETURN-1}}{(H_e, \ell_e, v), \text{return}_1 \cdot \Downarrow \text{ret}(H_e, \ell_e, v)}$$

### 2.3 Aborting Rules

$$\begin{array}{c}
 \text{RED-ERROR-EXPR}(e) \\
 \hline
 \sigma, e \Downarrow \text{err} \quad \mathbf{abort} \sigma \wedge \mathbf{-intercept}_e \sigma
 \end{array}
 \qquad
 \begin{array}{c}
 \text{RED-ERROR-STAT}(s) \\
 \hline
 \sigma, s \Downarrow \text{err} \quad \mathbf{abort} \sigma
 \end{array}$$
  

$$\begin{array}{c}
 \sigma = C[\text{err}] \\
 \hline
 \mathbf{abort} \sigma
 \end{array}
 \qquad
 \frac{}{\mathbf{intercept}_{@_3} \text{ret}(H_e, \ell_e, v)}$$