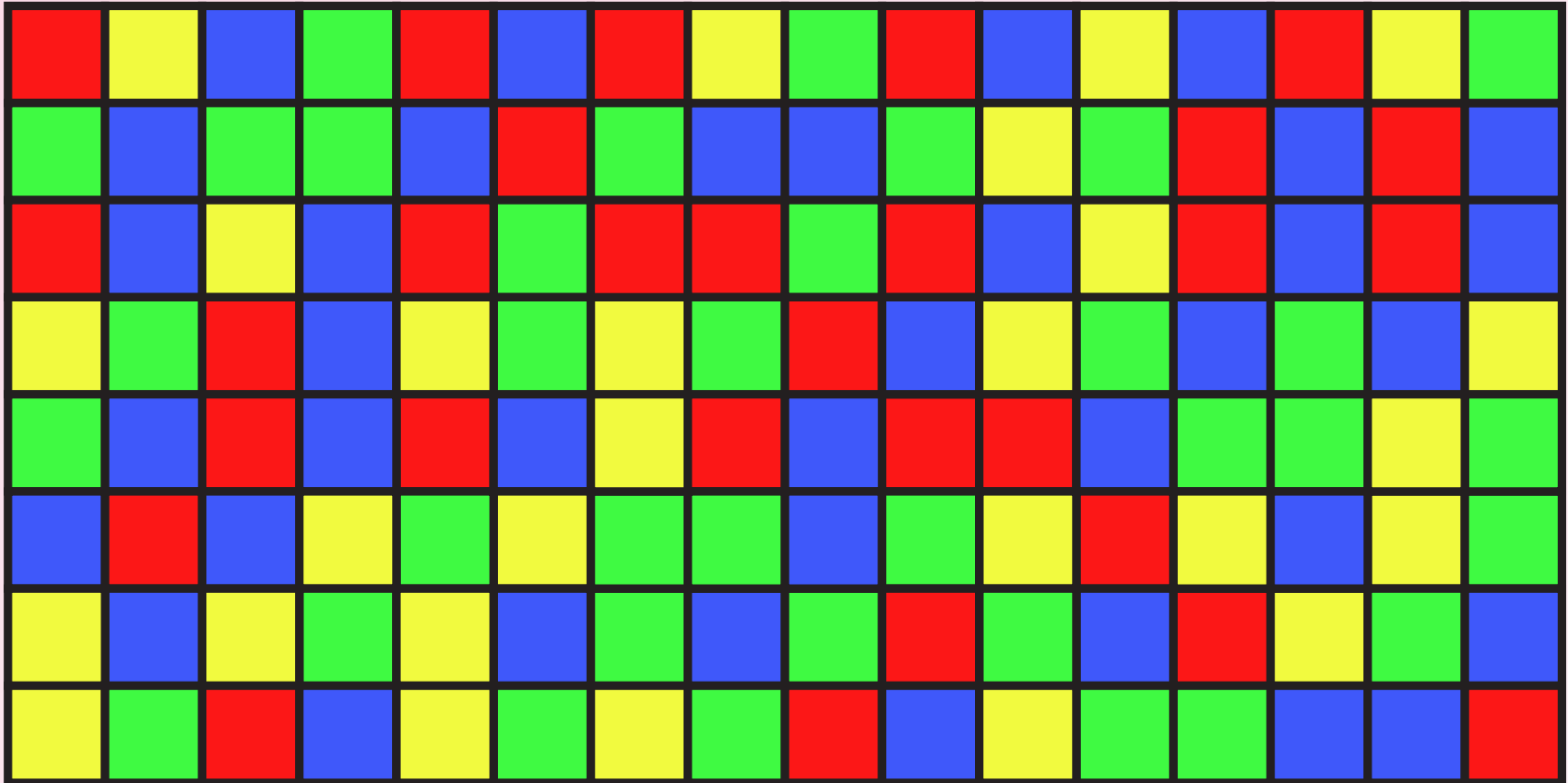


An algebraic geometric approach
to
multidimensional words

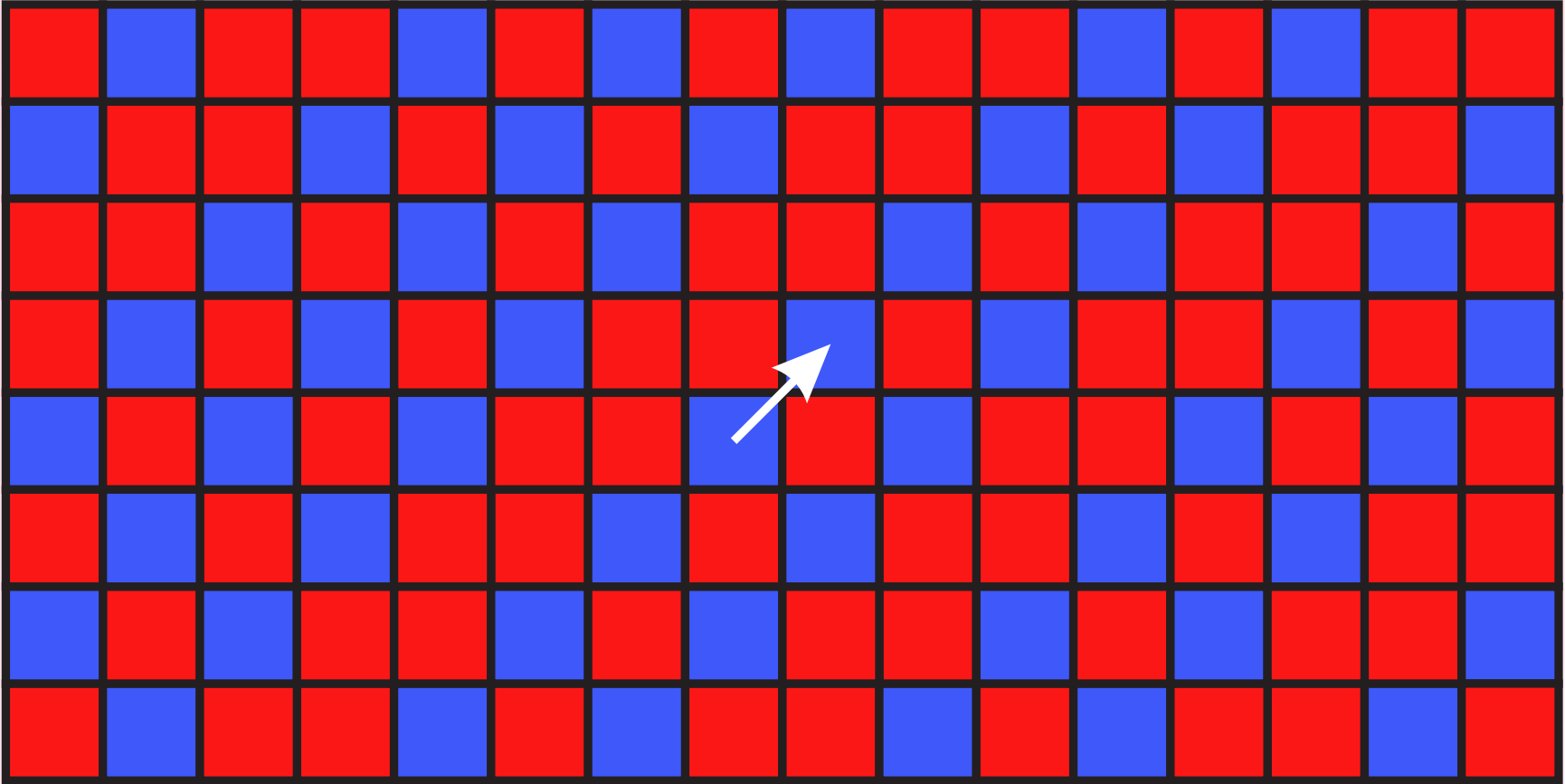
Jarkko Kari and Michal Szabados

Department of Mathematics and Statistics

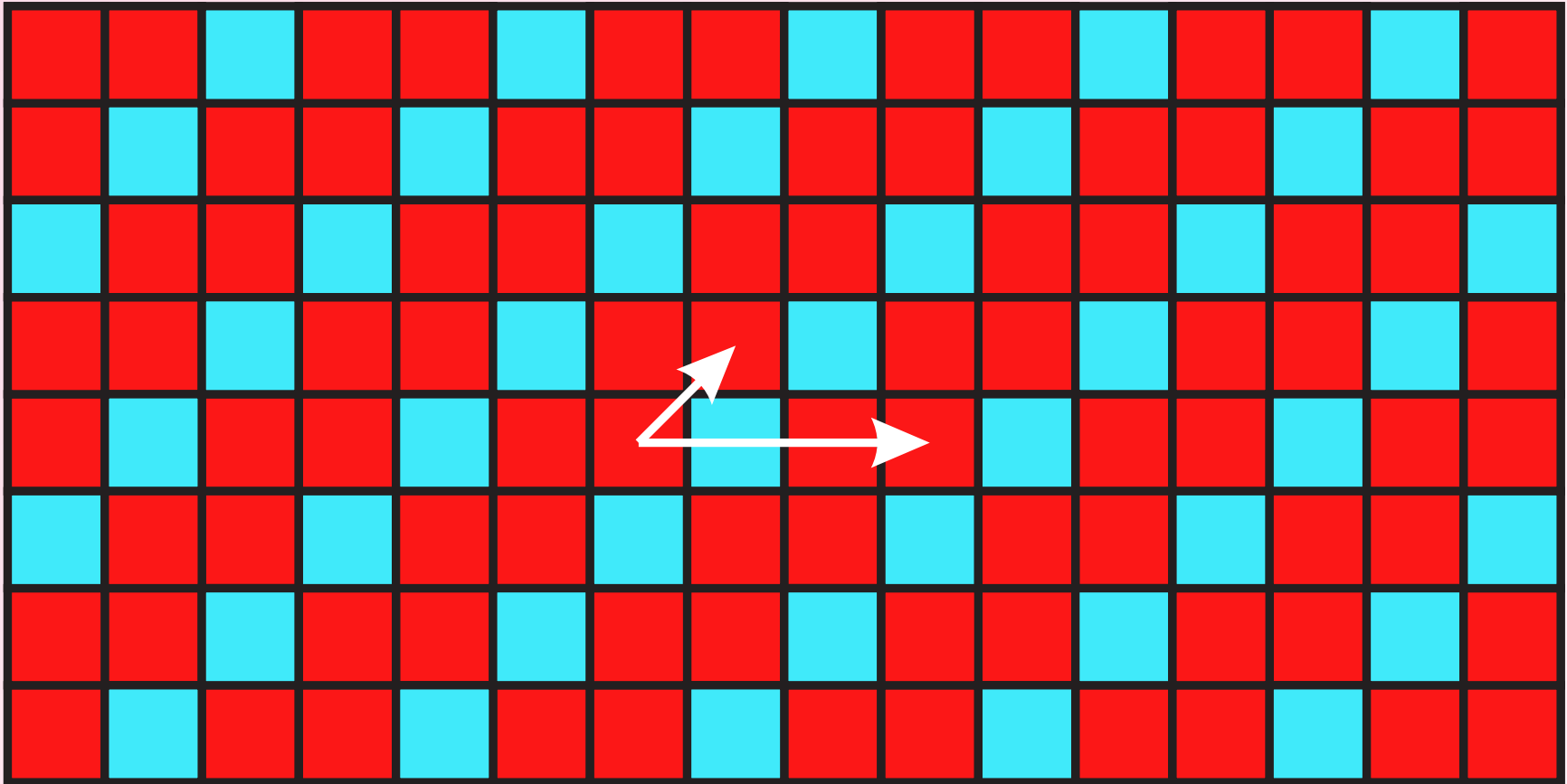
University of Turku, Finland



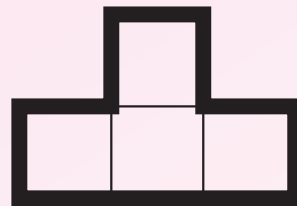
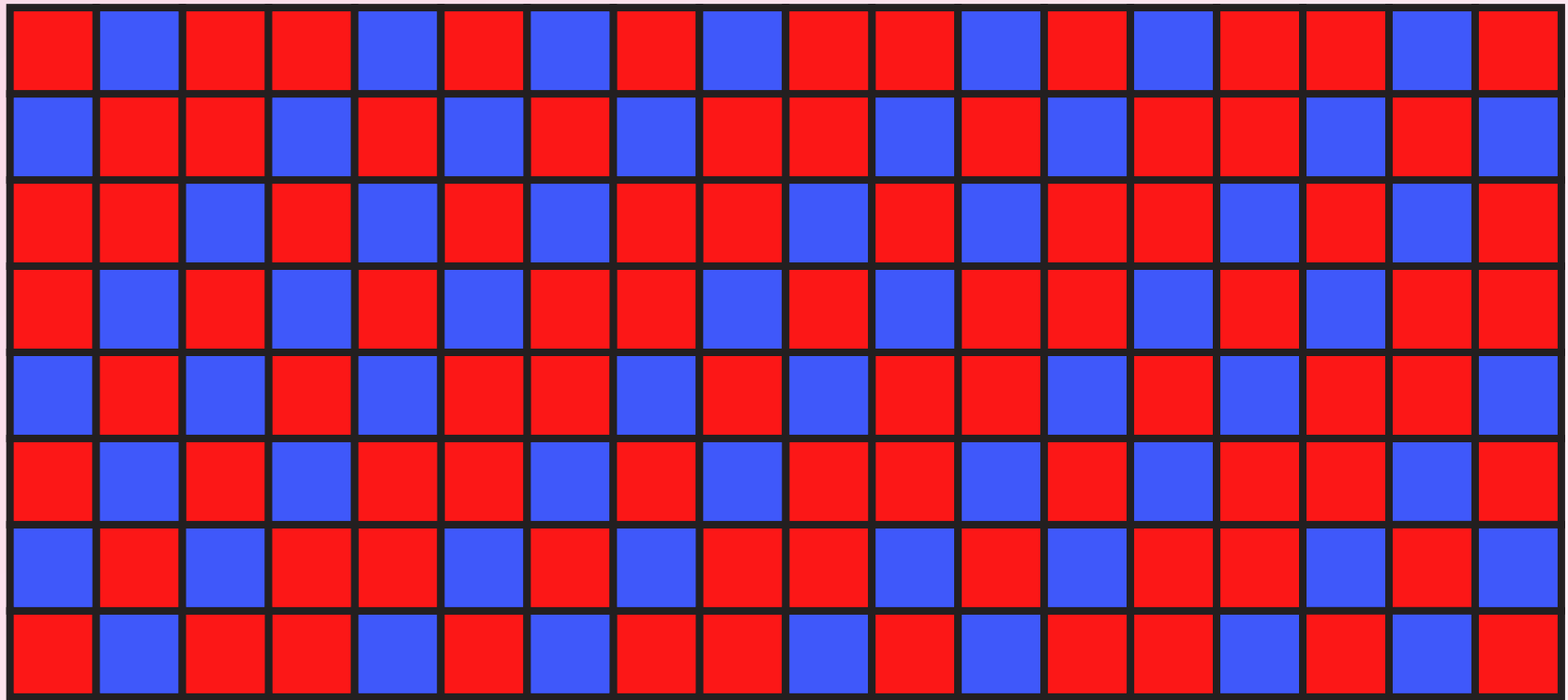
A d -dimensional **configuration** over set A is an element of $A^{\mathbb{Z}^d}$.



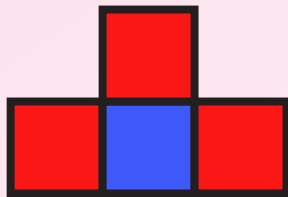
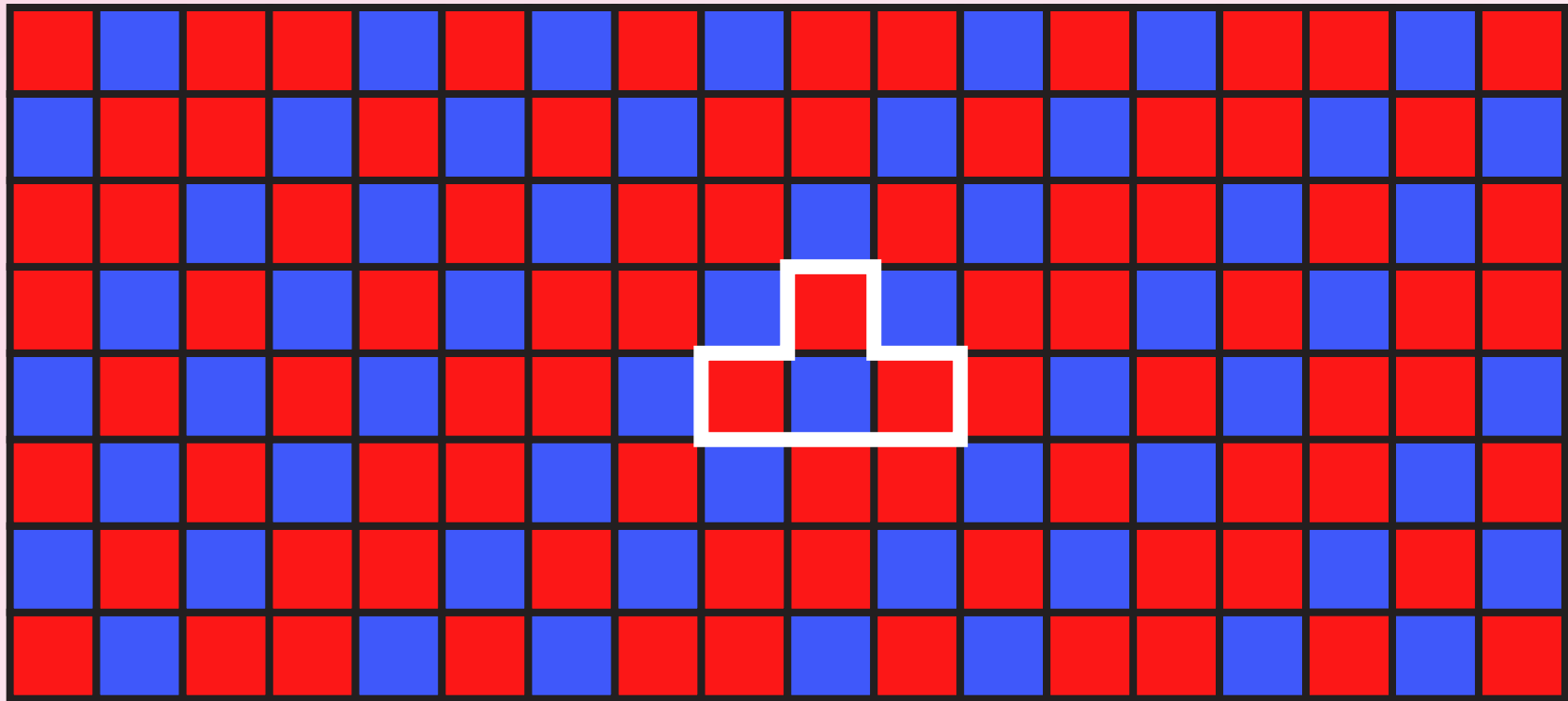
Configuration is **periodic** if it is invariant under a non-zero translation.



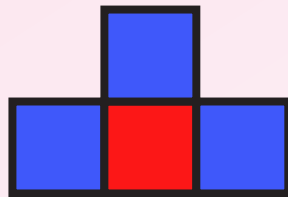
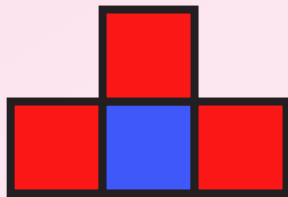
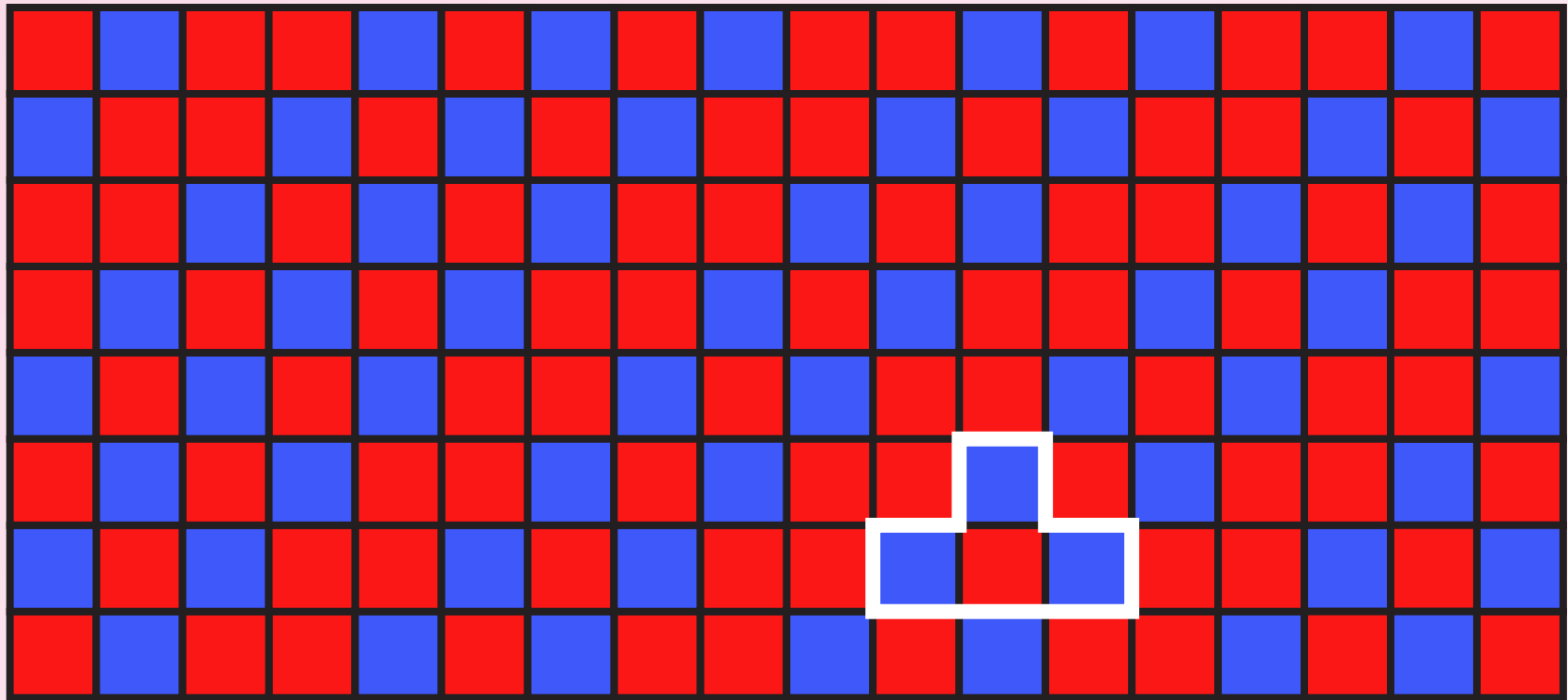
And a 2D configuration is **doubly periodic** if it is periodic in two directions (i.e. has horizontal and vertical periods).



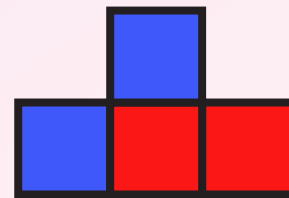
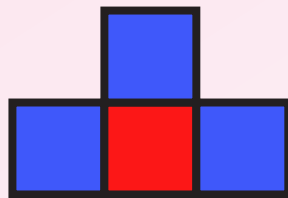
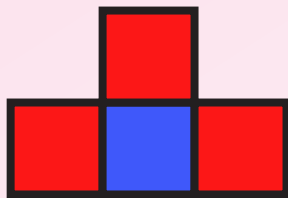
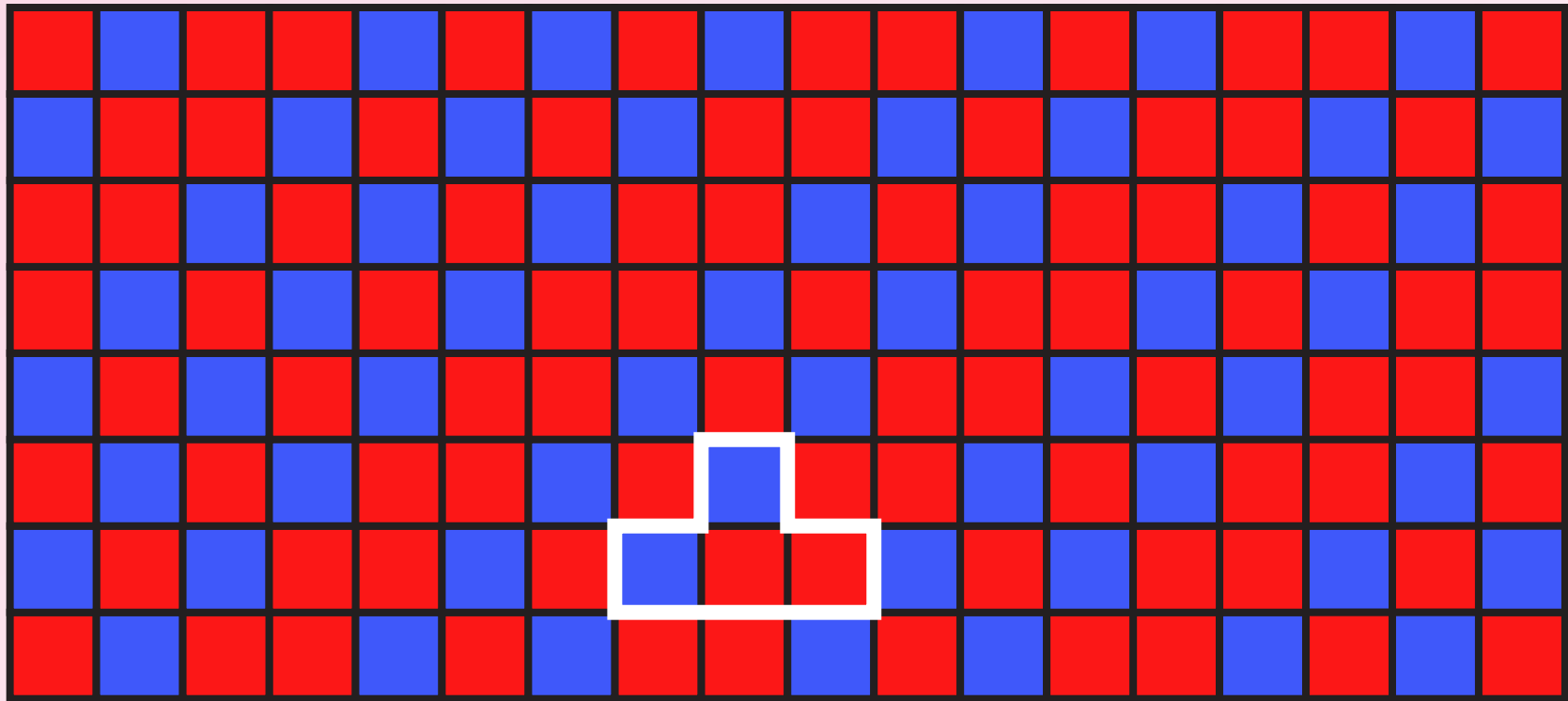
For finite **observation window** $D \subseteq \mathbb{Z}^d$ we call the patterns of shape D that appear in configuration c the **D -patterns** of c .



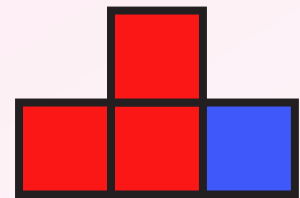
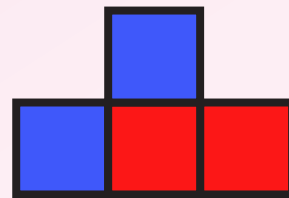
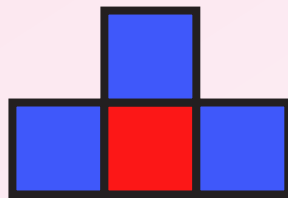
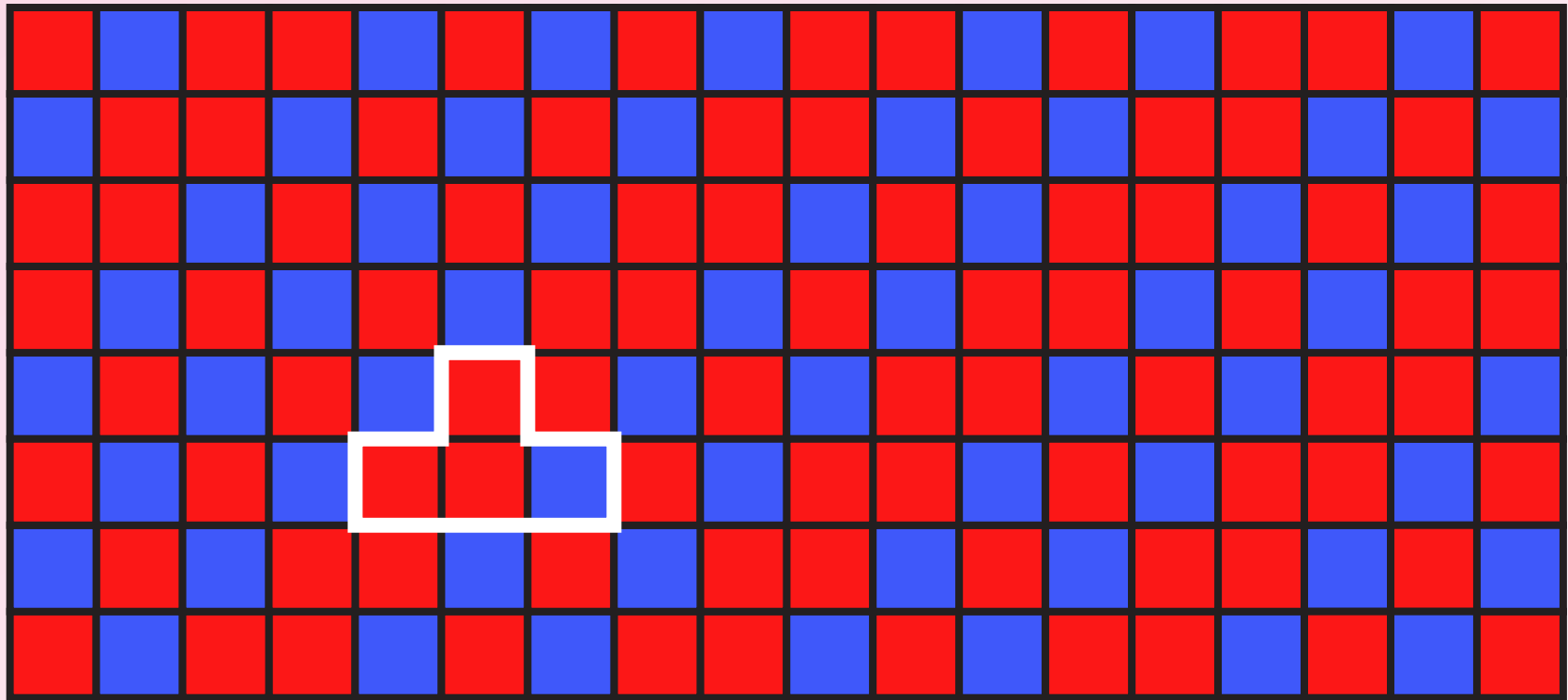
For finite **observation window** $D \subseteq \mathbb{Z}^d$ we call the patterns of shape D that appear in configuration c the **D -patterns** of c .



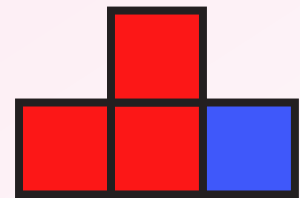
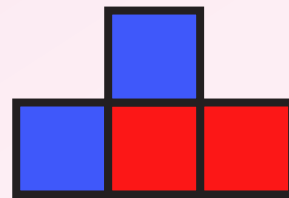
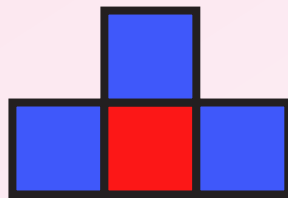
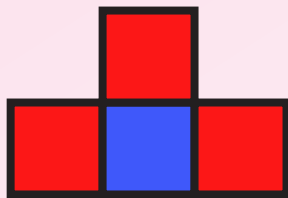
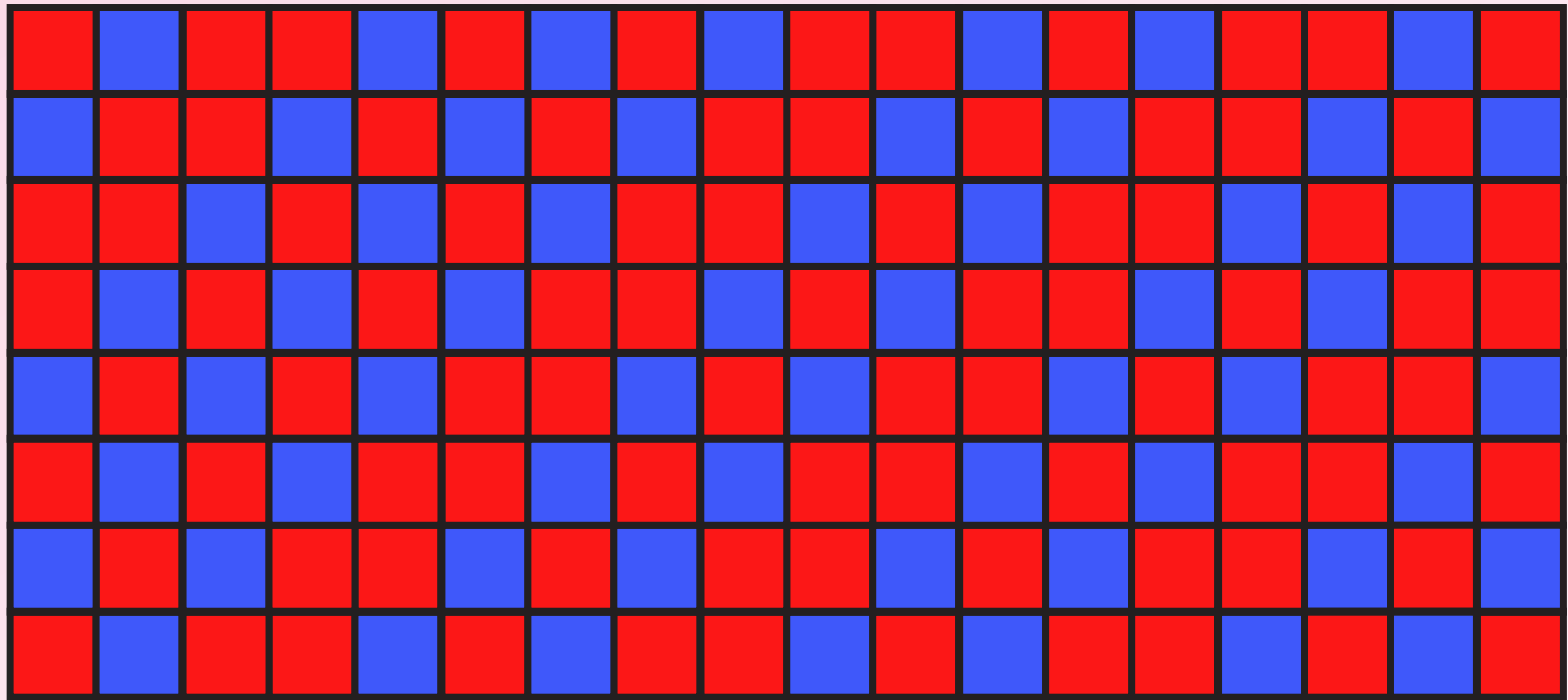
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For finite **observation window** $D \subseteq \mathbb{Z}^d$ we call the patterns of shape D that appear in configuration c the **D -patterns** of c .



The number of D -patterns is the pattern **complexity**

$$P(c, D) = \#\{\tau(c)_D \mid \tau \text{ is a translation}\}.$$

Local vs. global order

When does small pattern complexity imply periodicity ? Can a small pattern set prevent periodicity ?

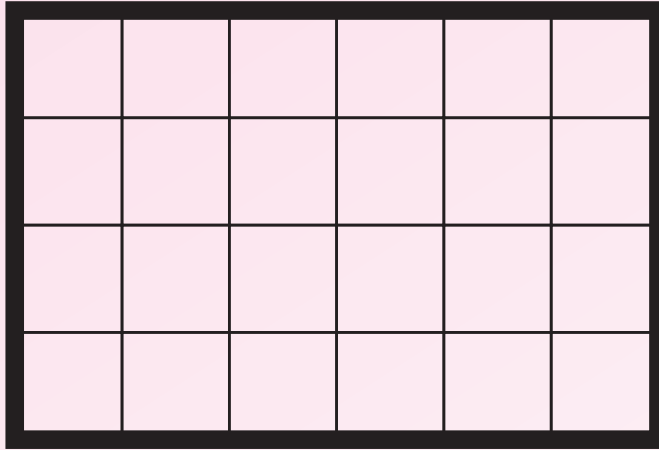
We are interested in the case when

$$P(c, D) \leq |D|$$

for some finite D .

Open problem 1: Nivat's conjecture

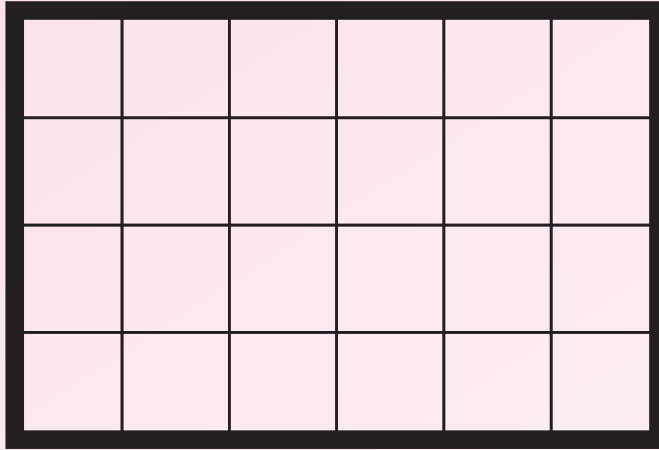
Consider $d = 2$ and rectangular D .



Conjecture (Nivat 1997) If $P(c, D) \leq |D|$ for some rectangle D then c is periodic.

Open problem 1: Nivat's conjecture

Consider $d = 2$ and rectangular D .



Conjecture (Nivat 1997) If $P(c, D) \leq |D|$ for some rectangle D then c is periodic.

This would extend the one-dimensional case $d = 1$:

Morse-Hedlund theorem: Let $c \in A^{\mathbb{Z}}$ and $n \in \mathbb{N}$. If c has at most n distinct subwords of length n then c is periodic.

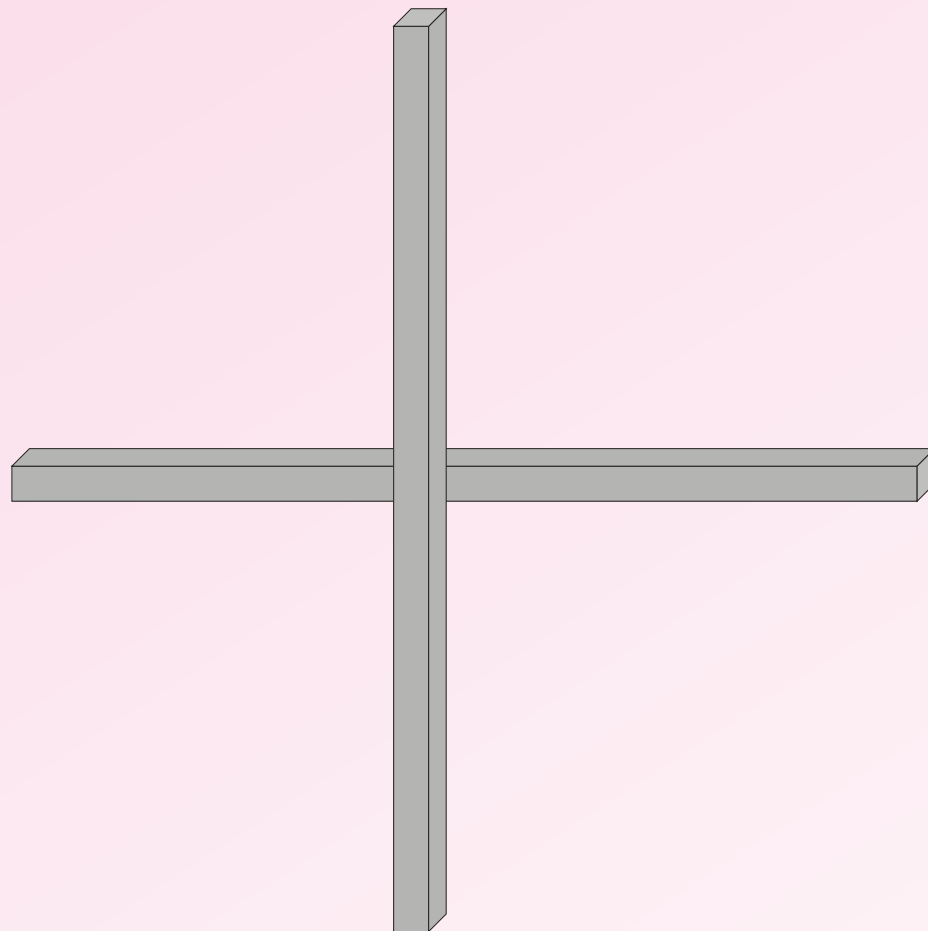
Best known bound in 2D:

Theorem (Cyr, Kra): If $P(c, D) \leq \frac{1}{2}|D|$ for some rectangle D then c is periodic.

Also narrow rectangles have been taken care of:

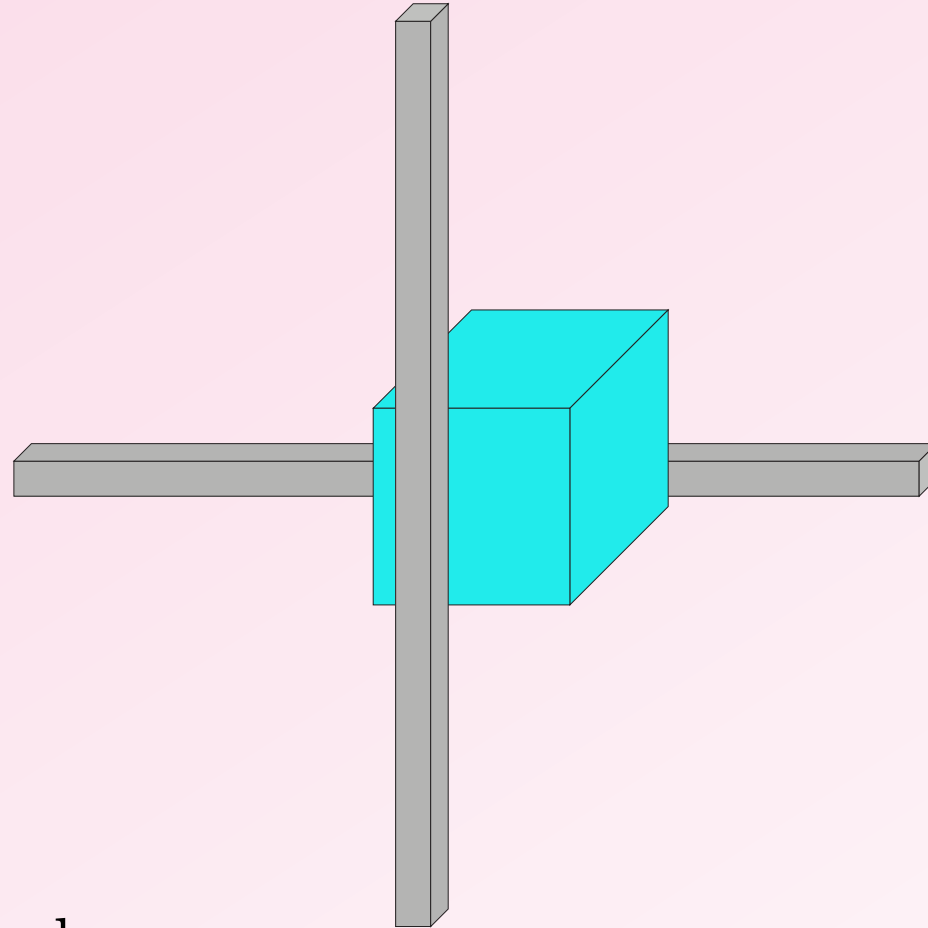
Theorem (Cyr, Kra): If D is a rectangle of height at most 3 and $P(c, D) \leq |D|$ then c is periodic.

In 3D and higher dimensional cases the conjecture is false



Non-periodic c

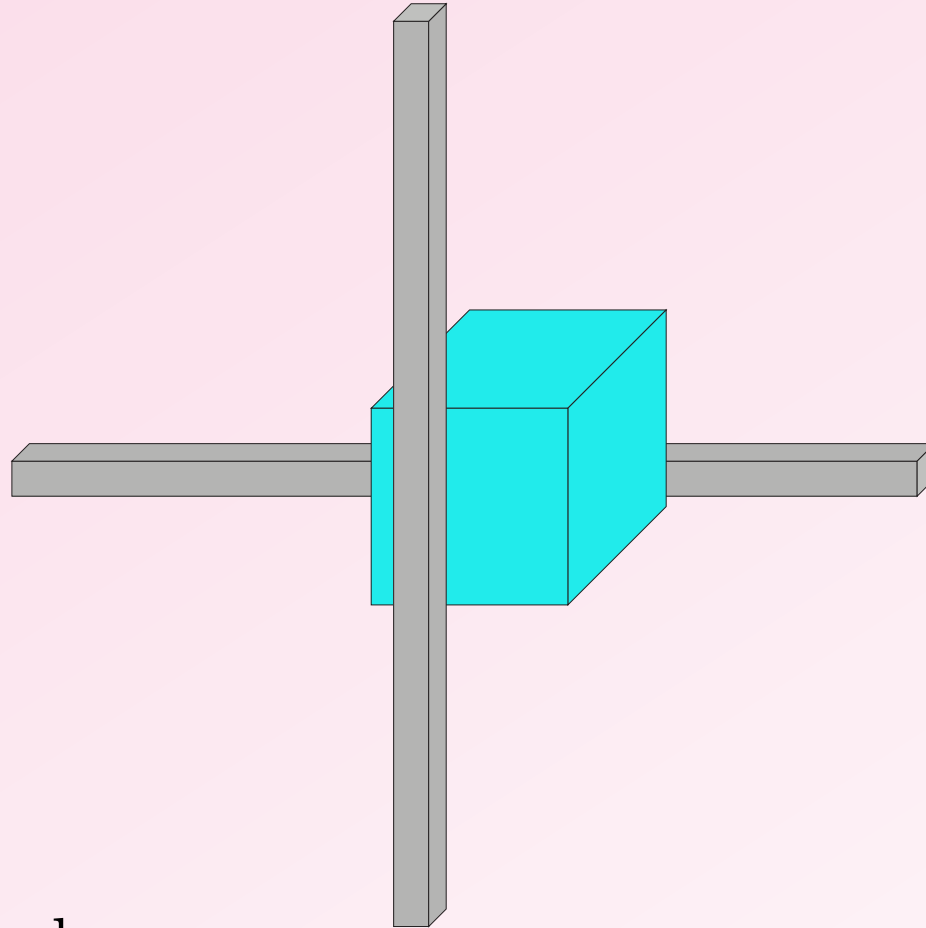
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Non-periodic c

D is $n \times n \times n$ cube

In 3D and higher dimensional cases the conjecture is false

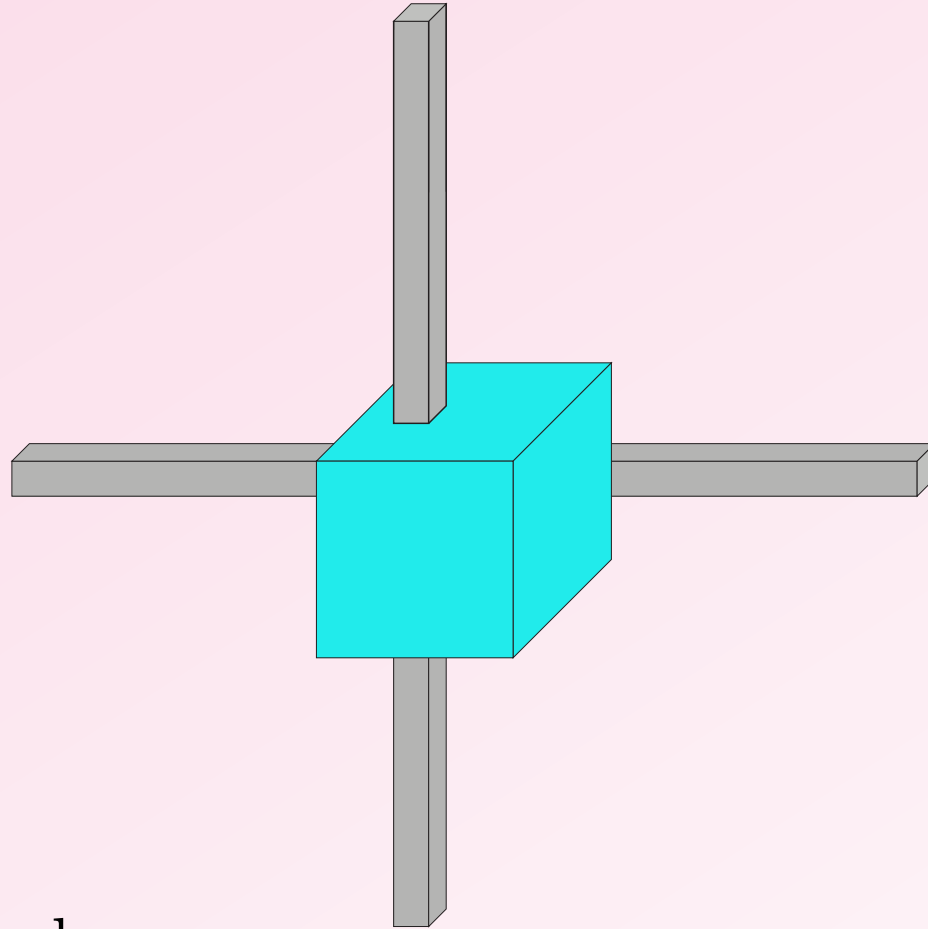


Non-periodic c

D is $n \times n \times n$ cube

$$P(c, D) = 1 + \dots$$

In 3D and higher dimensional cases the conjecture is false

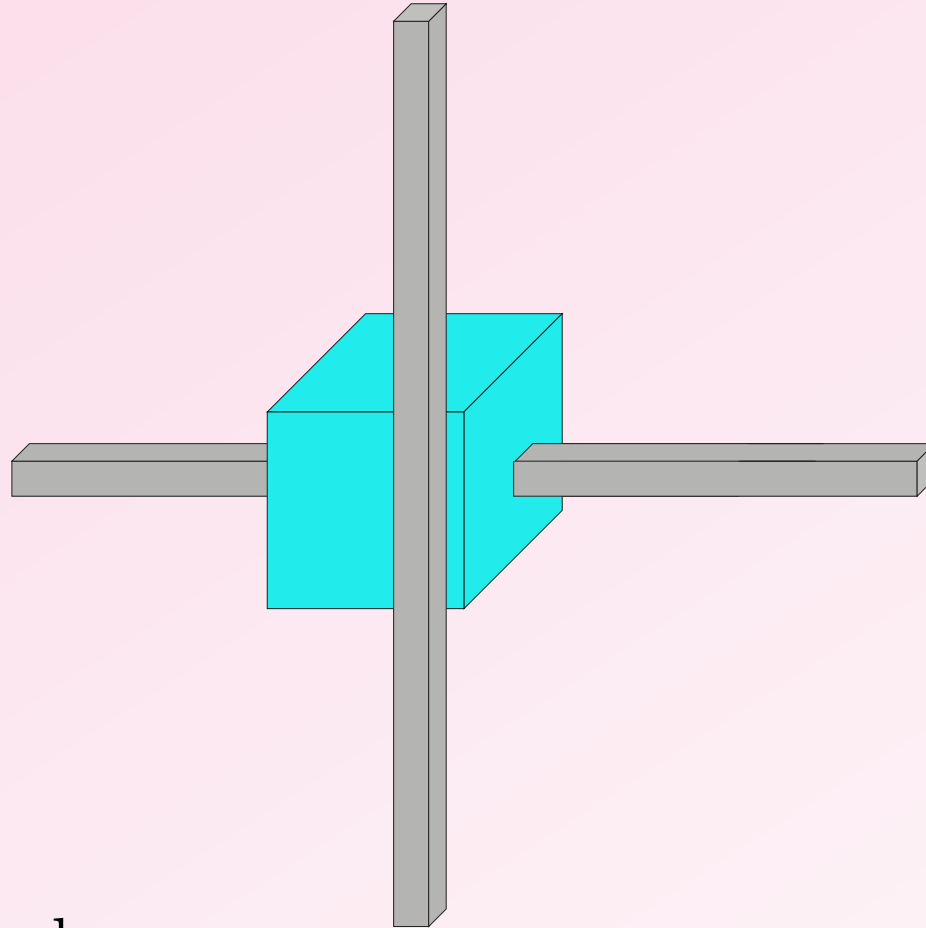


Non-periodic c

D is $n \times n \times n$ cube

$$P(c, D) = 1 + n^2 + \dots$$

In 3D and higher dimensional cases the conjecture is false

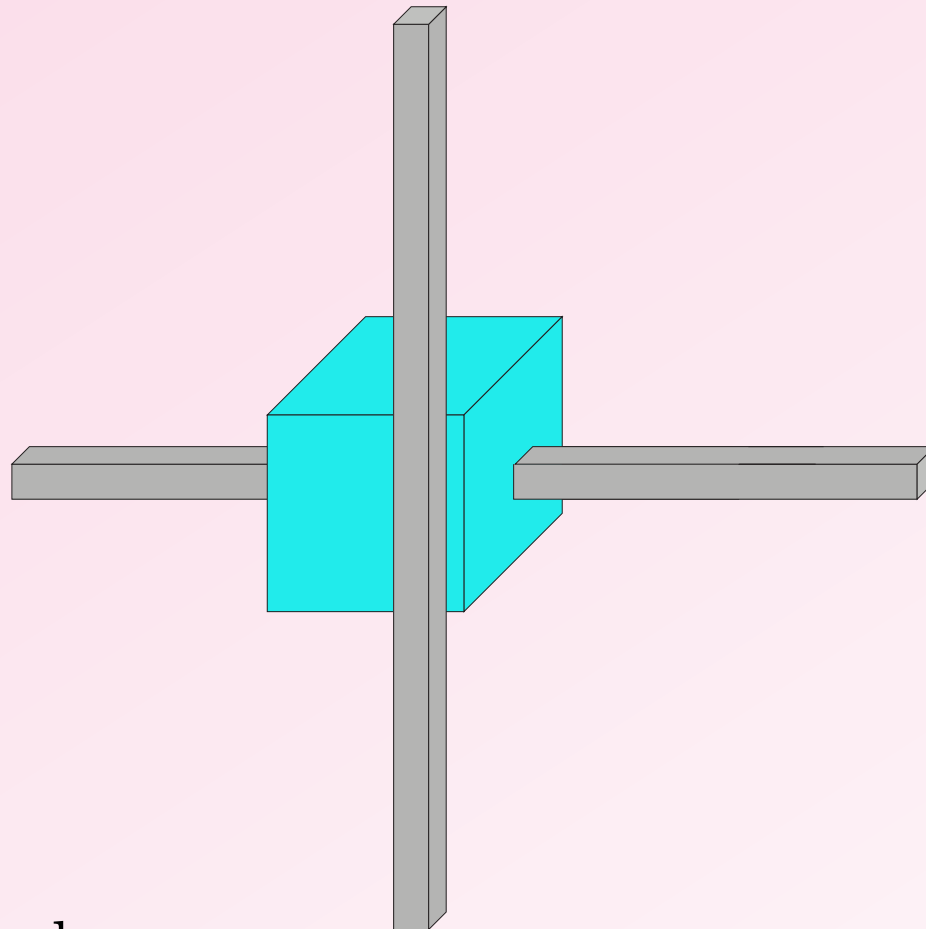


Non-periodic c

D is $n \times n \times n$ cube

$$P(c, D) = 1 + n^2 + n^2$$

In 3D and higher dimensional cases the conjecture is false



Non-periodic c

D is $n \times n \times n$ cube

$P(c, D) = 1 + n^2 + n^2 < n^3 = |D|$ for large n .

We can prove an asymptotic version in 2D:

Theorem (Kari, Szabados): If $P(c, D) \leq |D|$ for infinitely many different size rectangles D then c is periodic.

Open problem 2: Periodic tiling problem

Let $T \subseteq \mathbb{Z}^d$ be finite, and call it a **tile**. A **tiling** is any $C \subseteq \mathbb{Z}^d$ such that

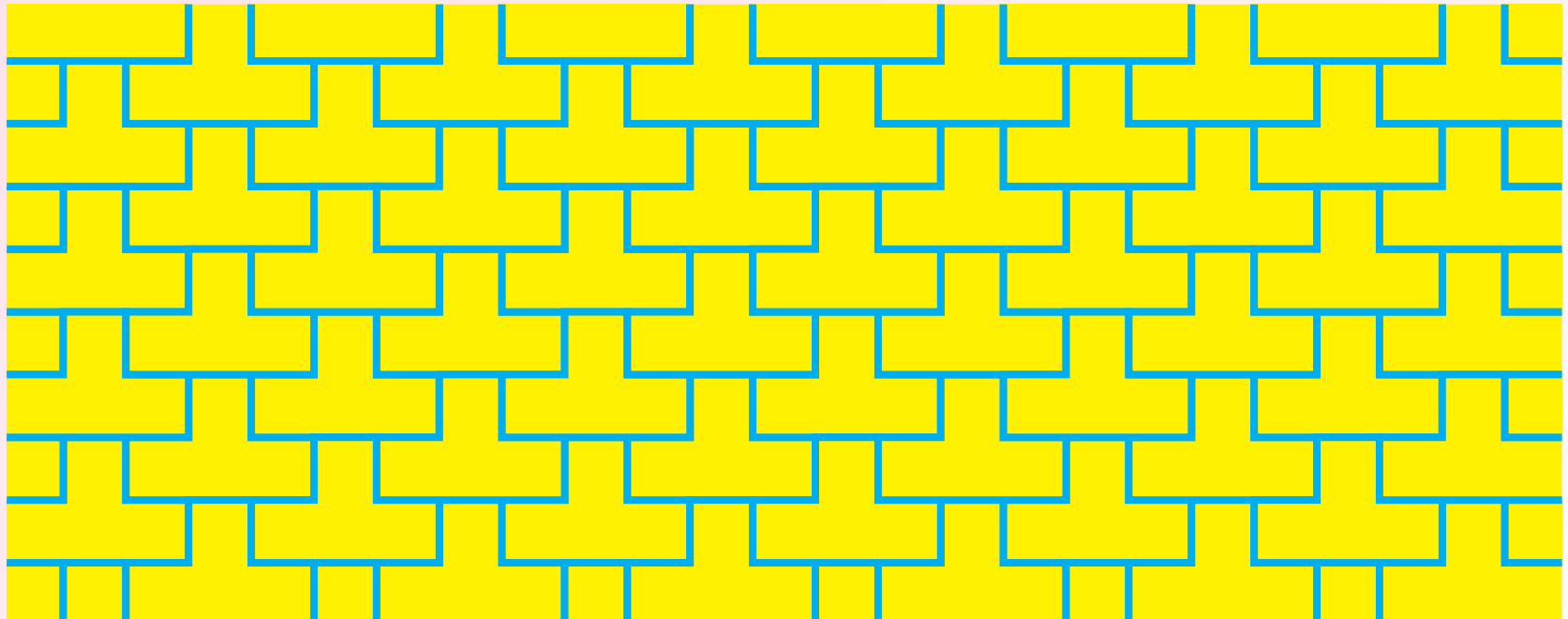
$$C \oplus T = \mathbb{Z}^d.$$

Open problem 2: Periodic tiling problem

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Graphical interpretation: C gives the positions where copies of T are placed to cover \mathbb{Z}^d without gaps or overlaps.

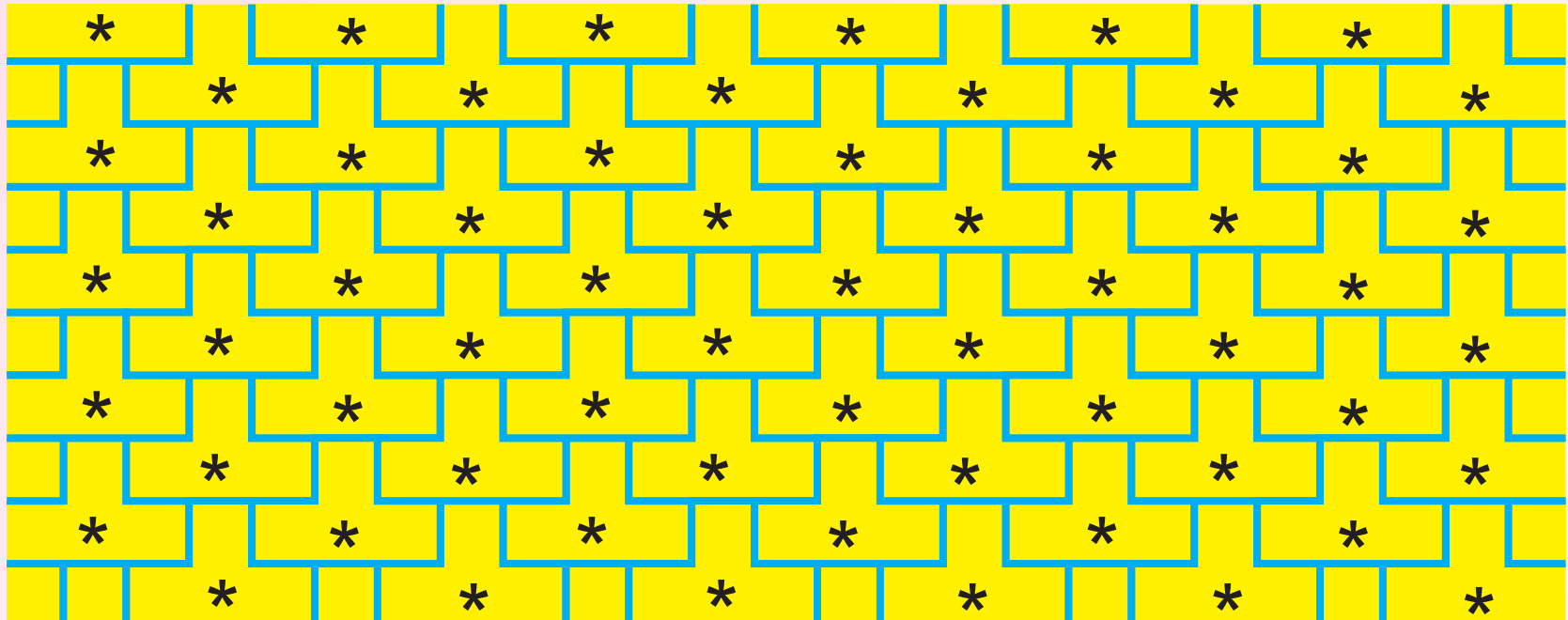


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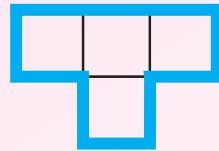


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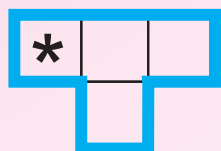
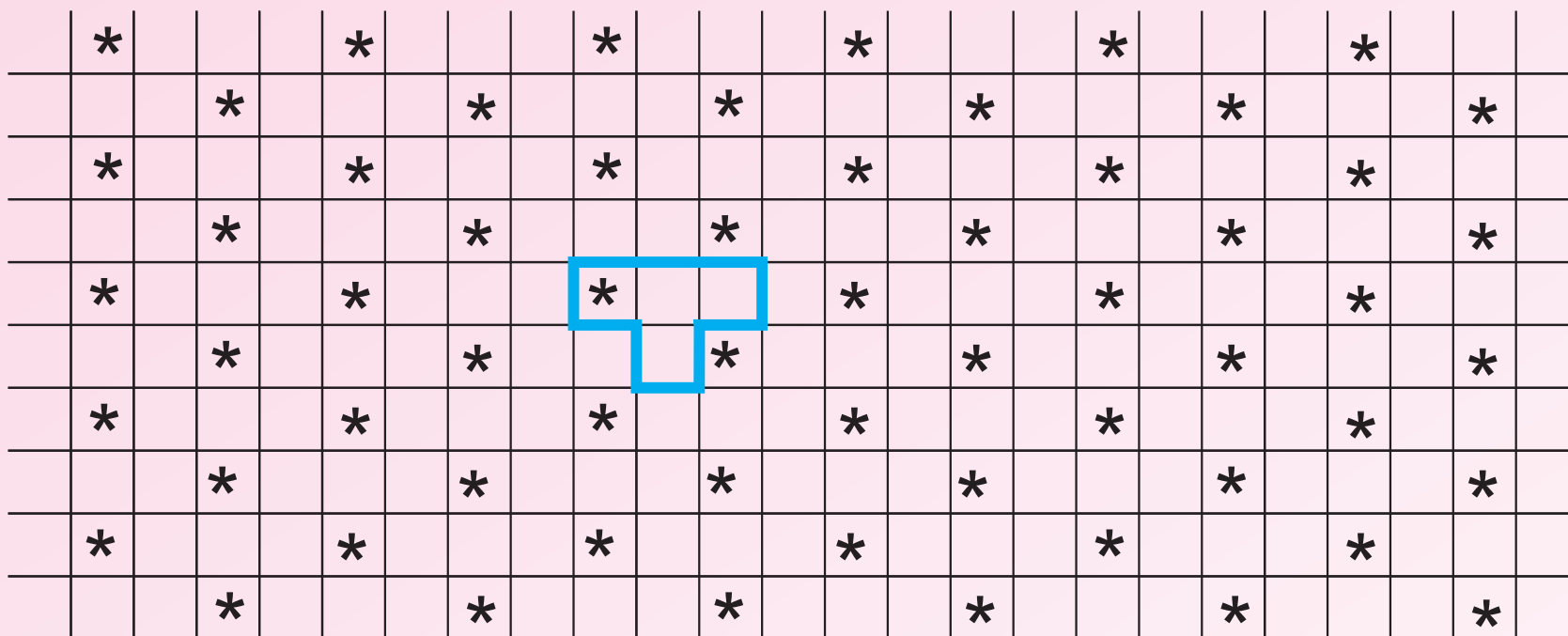
Interpret C as the binary configuration c with

$$c(i) = * \iff i \in C.$$

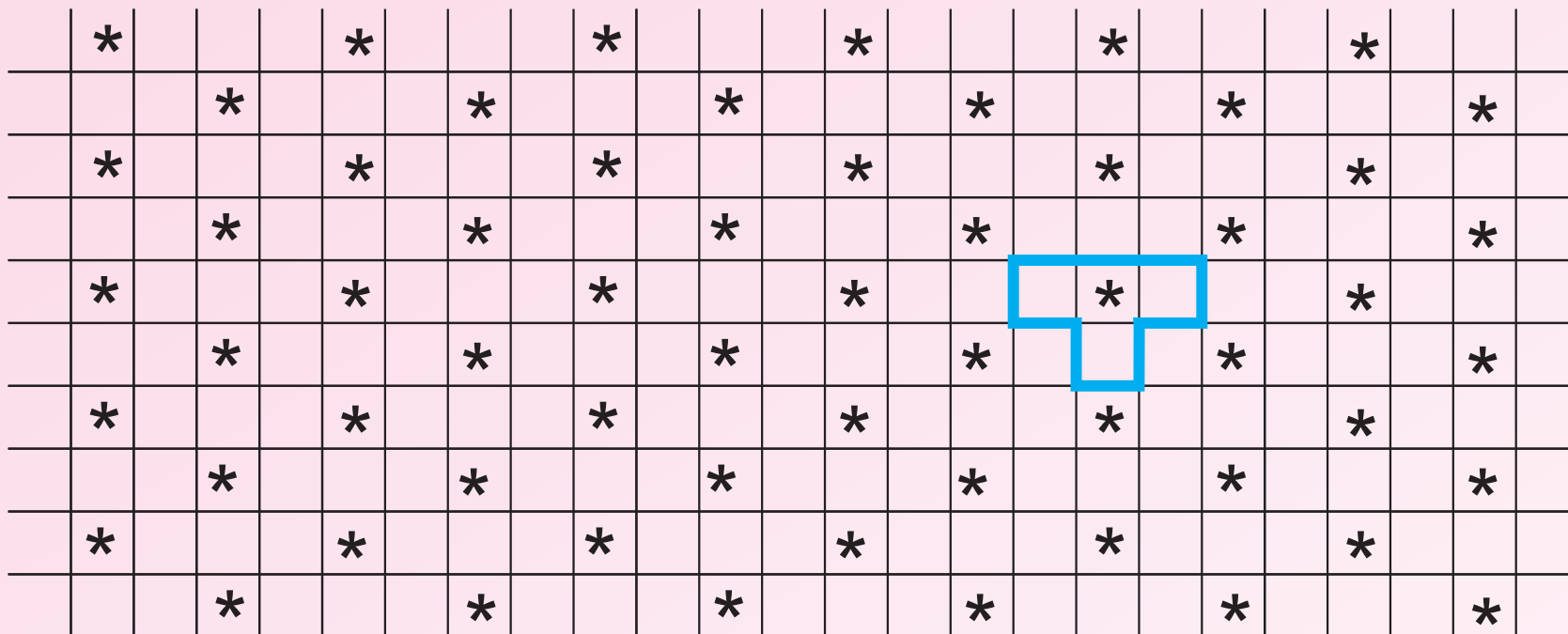
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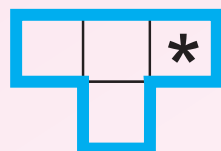
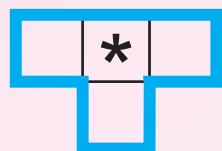
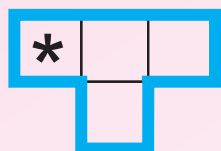
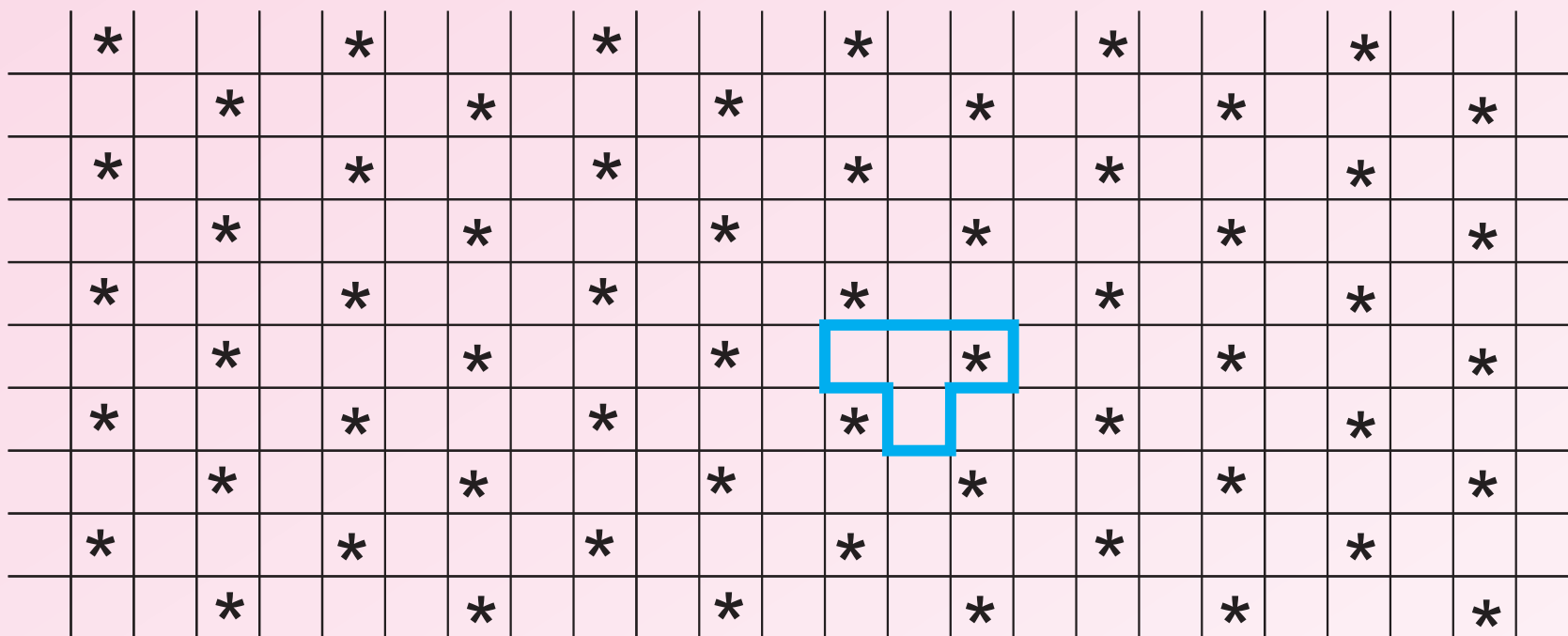
$(-T)$ -patterns of c contain exactly one symbol $*$.



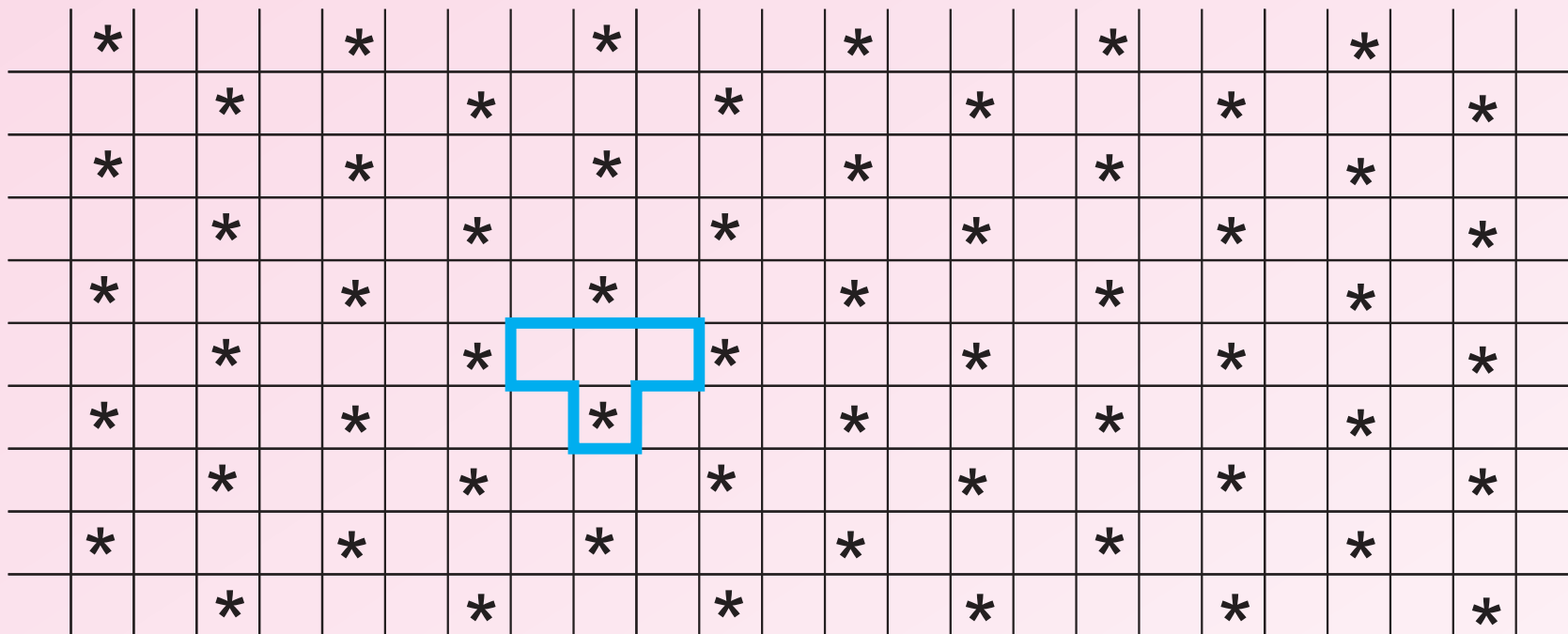
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$(-T)$ -patterns of c contain exactly one symbol $*$.

$$P(c, -T) = |T|$$

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$(-T)$ -patterns of c contain exactly one symbol $*$.

$$P(c, -T) = |T|$$

(Also $P(c, T) = |T|$ as any tiling for T is also a tiling for $-T$.)

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If X is the **SFT of all tilings** by T then

$$P(X, T) = |T|$$

where $P(X, T)$ is defined in the obvious way (=number of distinct T -patterns found in X).

Periodic tiling problem (Lagarias and Wang 1996): If T admits a tiling C , does it necessarily admit a periodic tiling ?

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Known results:

- Yes if $|T|$ is a prime number (Szegedy 1998).
- Yes in 2D if T is 4-connected (Beauquier and Nivat 1991).

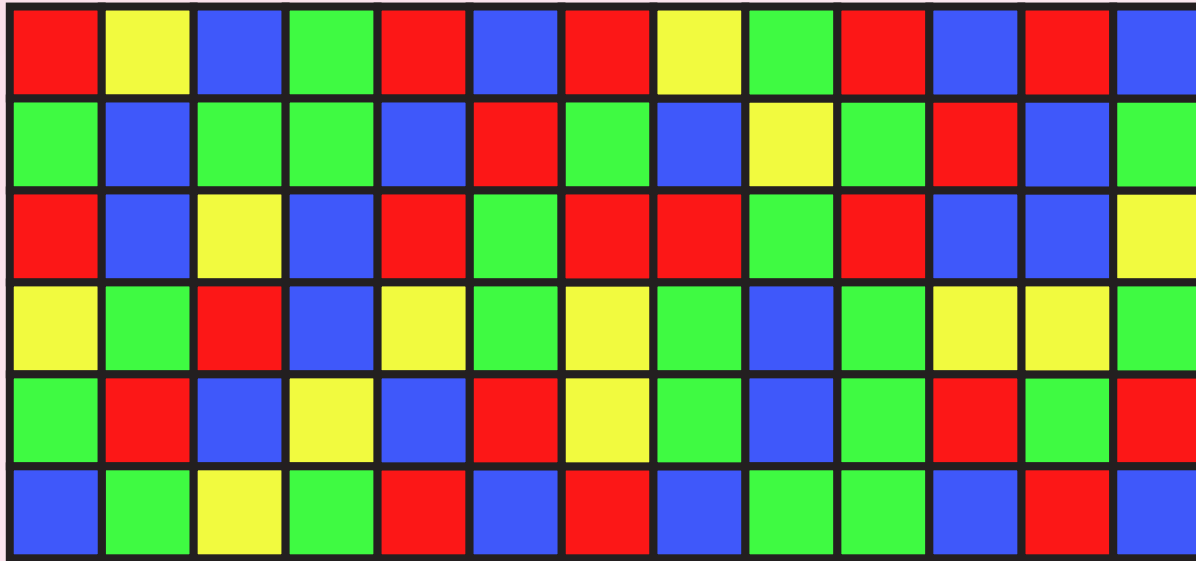
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Both the **Nivat's conjecture** and the **Periodic tiling problem** concern periodicity under complexity constraint $P(c, T) \leq |T|$.

We study configurations using algebra, so we first replace symbols by integers:



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1	2	3	4	1	3	1	2	4	1	3	1	3
4	3	4	4	3	1	4	3	2	4	1	3	4
1	3	2	3	1	4	1	1	4	1	3	3	2
2	4	1	3	2	4	2	4	3	4	2	2	4
4	1	3	2	3	1	2	4	3	4	1	4	1
3	4	2	4	1	3	1	3	4	4	3	1	3

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1	2	3	4	1	3	1	2	4	1	3	1	3
4	3	4	4	3	1	4	3	2	4	1	3	4
1	3	2	3	1	4	1	1	4	1	3	3	2
2	4	1	3	2	4	2	4	3	4	2	2	4
4	1	3	2	3	1	2	4	3	4	1	4	1
3	4	2	4	1	3	1	3	4	4	3	1	3

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4	3	4	4	3	1	4	3	2	4	1	3	4
1	3	2	3	1	4	1	1	4	1	3	3	2
2	4	1	3	2	4	2	4	3	4	2	2	4
4	1	3	2	3	1	2	4	3	4	1	4	1
3	4	2	4	1	3	1	3	4	4	3	1	3

D -patterns are viewed as $|D|$ -dimensional numerical vectors.

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1	2	3	4	1	3	1	2	4	1	3	1	3
4	3	4	4	3	1	4	3	2	4	1	3	4
1	3	2	3	1	4	1	1	4	1	3	3	2
2	4	1	3	2	4	2	4	3	4	2	2	4
4	1	3	2	3	1	2	4	3	4	1	4	1
3	4	2	4	1	3	1	3	4	4	3	1	3

D -patterns are viewed as $|D|$ -dimensional numerical vectors.

If $P(c, D) < |D|$ then there is an (integer) vector orthogonal to all D -patterns of c . **Filtering** c by this vector produces zero-configuration.

We study configurations using algebra, so we first replace symbols by integers:

1	2	3	4	1	3	1	2	4	1	3	1	3
4	3	4	4	3	1	4	3	2	4	1	3	4
1	3	2	3	1	4	1	1	4	1	3	3	2
2	4	1	3	2	4	2	4	3	4	2	2	4
4	1	3	2	3	1	2	4	3	4	1	4	1
3	4	2	4	1	3	1	3	4	4	3	1	3

If $P(c, D) \leq |D|$ we can add a suitable rational constant to c so that an orthogonal vector exists.

This is OK: we are free to choose the encoding of symbols as numbers freely.

We study configurations using algebra, so we first replace symbols by integers:

1	2	3	4	1	3	1	2	4	1	3	1	3
4	3	4	4	3	1	4	3	2	4	1	3	4
1	3	2	3	1	4	1	1	4	1	3	3	2
2	4	1	3	2	4	2	4	3	4	2	2	4
4	1	3	2	3	1	2	4	3	4	1	4	1
3	4	2	4	1	3	1	3	4	4	3	1	3

Conclusion: If $P(c, D) \leq |D|$ then symbols can be represented as integers in such a way that a non-trivial integer filter annihilates c .

We study configurations using algebra, so we first replace symbols by integers:

1	2	3	4	1	3	1	2	4	1	3	1	3
4	3	4	4	3	1	4	3	2	4	1	3	4
1	3	2	3	1	4	1	1	4	1	3	3	2
2	4	1	3	2	4	2	4	3	4	2	2	4
4	1	3	2	3	1	2	4	3	4	1	4	1
3	4	2	4	1	3	1	3	4	4	3	1	3

To use algebraic geometry, we next represent c as a power series (negative exponents included):

$$c \longleftrightarrow \sum_{(i_1, \dots, i_d) \in \mathbb{Z}^d} c(i_1, \dots, i_d) x_1^{i_1} \dots x_d^{i_d}$$

$$c \longleftrightarrow \sum_{(i_1, \dots, i_d) \in \mathbb{Z}^d} c(i_1, \dots, i_d) x_1^{i_1} \dots x_d^{i_d}$$

Notations:

- $X = (x_1, \dots, x_d)$
- For $I = (i_1, \dots, i_d) \in \mathbb{Z}^d$ we denote by

$$X^I = x_1^{i_1} \dots x_d^{i_d}$$

the monomial that represents cell I .

- For $n \in \mathbb{Z}$ we use $X^n = X^{(n, \dots, n)} = x_1^n \dots x_d^n$.

$$c \longleftrightarrow \sum_{(i_1, \dots, i_d) \in \mathbb{Z}^d} c(i_1, \dots, i_d) x_1^{i_1} \dots x_d^{i_d} = \sum_{I \in \mathbb{Z}^d} c(I) X^I$$

Notations:

- $X = (x_1, \dots, x_d)$
- For $I = (i_1, \dots, i_d) \in \mathbb{Z}^d$ we denote by

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More notations:

Let R be a ring. (Usually \mathbb{Z} or \mathbb{C} .)

- We denote $Z^I = z_1^{i_1} \dots z_d^{i_d}$ for any

$$Z = (z_1, \dots, z_d) \in R^d,$$

$$I = (i_1, \dots, i_d) \in \mathbb{Z}^d.$$

- $R[X]$ is the set of polynomials over ring R .
- $R[X^{\pm 1}]$ is the set of Laurent polynomials over ring R .
- $R[[X^{\pm 1}]]$ is the set of power series (negative exponents included) over ring R .

A filter is a Laurent polynomial $f(X)$, and filtering corresponds to multiplying a power series with $f(X)$.

We say that $f(X)$ **annihilates** $c(X)$ if $f(X)c(X) = 0$. Zero polynomial $f(X) = 0$ annihilates every configuration – it is the **trivial annihilator**.

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Our setup is

$$c(X) \in \mathbb{Z}[[X^{\pm 1}]]$$

that has non-trivial annihilators

$$f(X) \in \mathbb{Z}[X^{\pm 1}].$$

Monomial X^I represents the **translation** by vector $I \in \mathbb{Z}^d$. If $f(X)$ annihilates $c(X)$ then so does $X^I \cdot f(X)$.

For any Laurent polynomial $f(X)$ there exists I such that $X^I \cdot f(X)$ is a polynomial.

\implies our $c(X)$ has non-trivial **annihilating polynomials**.

Define

$$\text{Ann}(c) = \{f(X) \in \mathbb{C}[X] \mid f(X)c(X) = 0\}.$$

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Remarks:

- We consider polynomials (not Laurent polynomials!) so that we can directly rely on polynomial algebra.
- We allow complex coefficients because we need algebraically closed field to apply Hilbert's Nullstellensatz.
- $\text{Ann}(c)$ is indeed an ideal of the polynomial ring $\mathbb{C}[X]$.

In the following we prove that $\text{Ann}(c)$ contains polynomials of “simple form”, using

Nullstellensatz (Hilbert): Let $g(X)$ be a polynomial.

Suppose that $g(Z) = 0$ for all Z in the **variety**

$$\{Z \in \mathbb{C}^d \mid f(Z) = 0 \text{ for all } f \in \text{Ann}(c) \}.$$

Then $g^k \in \text{Ann}(c)$ for some $k \in \mathbb{N}$.

In the following, always, $c(X)$ is a finitary, integral power series and

$$f(X) = \sum_{I \in \mathcal{I}} a_I X^I$$

is its non-trivial integral annihilator polynomial, $a_I \neq 0$ for all $I \in \mathcal{I}$.

Lemma 1: There is an integer M such that $f(X^n) \in \text{Ann}(c)$ for every $n \in \mathbb{N}$ that is relatively prime with M .

Number M only depends on $c(X)$ and the coefficients of $f(X)$.
(Same M works for all polynomials with the same coefficients.)

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$$f(X) \begin{array}{|c|c|c|} \hline & a & \\ \hline b & c & d \\ \hline \end{array}$$

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$$f(X^2) \quad \begin{array}{ccc} & \boxed{a} & \\ \boxed{b} & \boxed{c} & \boxed{d} \end{array}$$

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$$f(X^3) \quad \begin{array}{c} \boxed{a} \\ \boxed{b} \quad \boxed{c} \quad \boxed{d} \end{array}$$

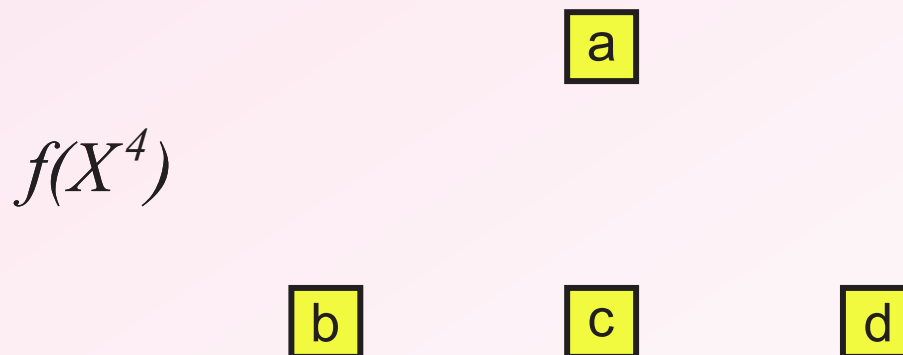
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Lemma 1: There is an integer M such that $f(X^n) \in \text{Ann}(c)$ for every $n \in \mathbb{N}$ that is relatively prime with M .

Proof: First consider the case that $n = p$ is a prime. Then

$$f(X)^p \equiv f(X^p) \pmod{p\mathbb{Z}[X]}$$

so because $f(X)c(X) = 0$ we have

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The coefficients in $f(X^p)c(X)$ are bounded in absolute value by

$$m = \sum_{I \in \mathcal{I}} |a_I| c_{\max}$$

where c_{\max} is the maximum absolute value of the coefficients in c .

If $p > m$ then $f(X^p)c(X) = 0$. We choose $M = m!$ so any number relatively prime with M is $> m$.

Lemma 1: There is an integer M such that $f(X^n) \in \text{Ann}(c)$ for every $n \in \mathbb{N}$ that is relatively prime with M .

Proof: General case $n = p_1 \dots p_k$ where primes p_k are relatively prime with M .

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Proof: General case $n = p_1 \dots p_k$ where primes p_k are relatively prime with M .

Using the prime case we have

$$f(X^{p_1}) \in \text{Ann}(c) \text{ so}$$

$$f(X^{p_1 p_2}) \in \text{Ann}(c) \text{ so}$$

...

$$f(X^{p_1 p_2 \dots p_k}) \in \text{Ann}(c).$$

Note that the coefficients of all $f(X^n)$ are the same as in $f(X)$ so the same M can be used everywhere. □

Lemma 1: There is an integer M such that $f(X^n) \in \text{Ann}(c)$ for every $n \in \mathbb{N}$ that is relatively prime with M .

In particular, $f(X^{1+iM})$ are in $\text{Ann}(c)$ for $i = 0, 1, 2, \dots$.

Let $Z \in \mathbb{C}^d$ be a common zero of $\text{Ann}(c)$. Then

$$f(Z^{1+iM}) = 0 \text{ for all } i = 0, 1, 2, \dots$$

In the following we show that $g(Z) = 0$ where

$$g(X) = X^1 \prod_{\substack{I, J \in \mathcal{I} \\ I \neq J}} (X^{MI} - X^{MJ}).$$

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Case 1: Some coordinate z_i of $Z = (z_1, \dots, z_d)$ is zero.

Then $Z^{(1, \dots, 1)} = 0$. Because Z is a zero of the monomial X^1 ,

$$g(Z) = 0.$$

Let $Z \in \mathbb{C}^d$ be a common zero of $\text{Ann}(c)$. Then

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Case 2: Assume then that all coordinates $z_i \neq 0$.

We partition the monomials of $f(X)$ according to the value of X^I for $X = Z^M$. For any $\alpha \in \mathbb{C}$ denote

$$\mathcal{I}_\alpha = \{I \in \mathcal{I} \mid Z^{MI} = \alpha\}.$$

Then \mathcal{I} is partitioned into some $\mathcal{I}_{\alpha_1}, \dots, \mathcal{I}_{\alpha_n}$.

$$f(X) = \sum_{I \in \mathcal{I}} a_I X^I$$

$$f(Z^{1+iM}) = 0, \quad \text{for all } i = 0, 1, 2, \dots$$

$$\mathcal{I}_\alpha = \{I \in \mathcal{I} \mid Z^{MI} = \alpha\}$$

$$\begin{aligned} f(X) &= \sum_{I \in \mathcal{I}} a_I X^I \\ f(Z^{1+iM}) &= 0, \quad \text{for all } i = 0, 1, 2, \dots \\ \mathcal{I}_\alpha &= \{I \in \mathcal{I} \mid Z^{MI} = \alpha\} \end{aligned}$$

Denote

$$f_\alpha(X) = \sum_{I \in \mathcal{I}_\alpha} a_I X^I.$$

Then

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Lemma 2. $f_\alpha(Z) = 0$ for all $\alpha = \alpha_1, \alpha_2, \dots, \alpha_n$.

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In other words, for all i vectors

$(f_{\alpha_1}(Z), \dots, f_{\alpha_n}(Z))$ and $(\alpha_1^i, \alpha_2^i, \dots, \alpha_n^i)$

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$$(f_{\alpha_1}(Z), \dots, f_{\alpha_n}(Z)) \quad \text{and} \quad (\alpha_1^i, \alpha_2^i, \dots, \alpha_n^i)$$

are orthogonal. This implies that all $f_{\alpha_i}(Z) = 0$. □

$$f(X) = \sum_{I \in \mathcal{I}} a_I X^I$$

$$f(Z^{1+iM}) = 0, \quad \text{for all } i = 0, 1, 2, \dots$$

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\end{aligned}$$

Because all coordinates z_i of Z are non-zero, $Z^I \neq 0$ for all I .

If some \mathcal{I}_α would contain exactly one $I \in \mathcal{I}$ then

$$0 = f_\alpha(Z) = a_I Z^I \neq 0.$$

So any non-empty \mathcal{I}_α contains at least two distinct $I, J \in \mathcal{I}$.

But $Z^{MI} = \alpha = Z^{MJ}$, so Z is a zero of $X^{MI} - X^{MJ}$.

We have proved the following: All elements of the variety

$$\{Z \in \mathbb{C}^d \mid f(Z) = 0 \text{ for all } f \in \text{Ann}(c) \}$$

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Nullstellensatz $\implies g(X)^n \in \text{Ann}(c)$ for some $n \in \mathbb{N}$.

Dividing $g(X)^n$ by a suitable monomial gives:

Theorem. Finitary, integral $c(X)$ that has a non-trivial annihilator is annihilated by a Laurent polynomial of the form

$$(1 - X^{I_1})(1 - X^{I_2}) \dots (1 - X^{I_k}).$$

$$\text{Annihilator: } (1 - X^{I_1})(1 - X^{I_2}) \dots (1 - X^{I_k})$$

Binomials $(1 - X^I)$ correspond to **difference operators** that subtract from a configuration its own I -translation.

The theorem states that configuration $c(X)$ can be annihilated by a sequence of difference operations.

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Binomials $(1 - X^I)$ correspond to **difference operators** that subtract from a configuration its own I -translation.

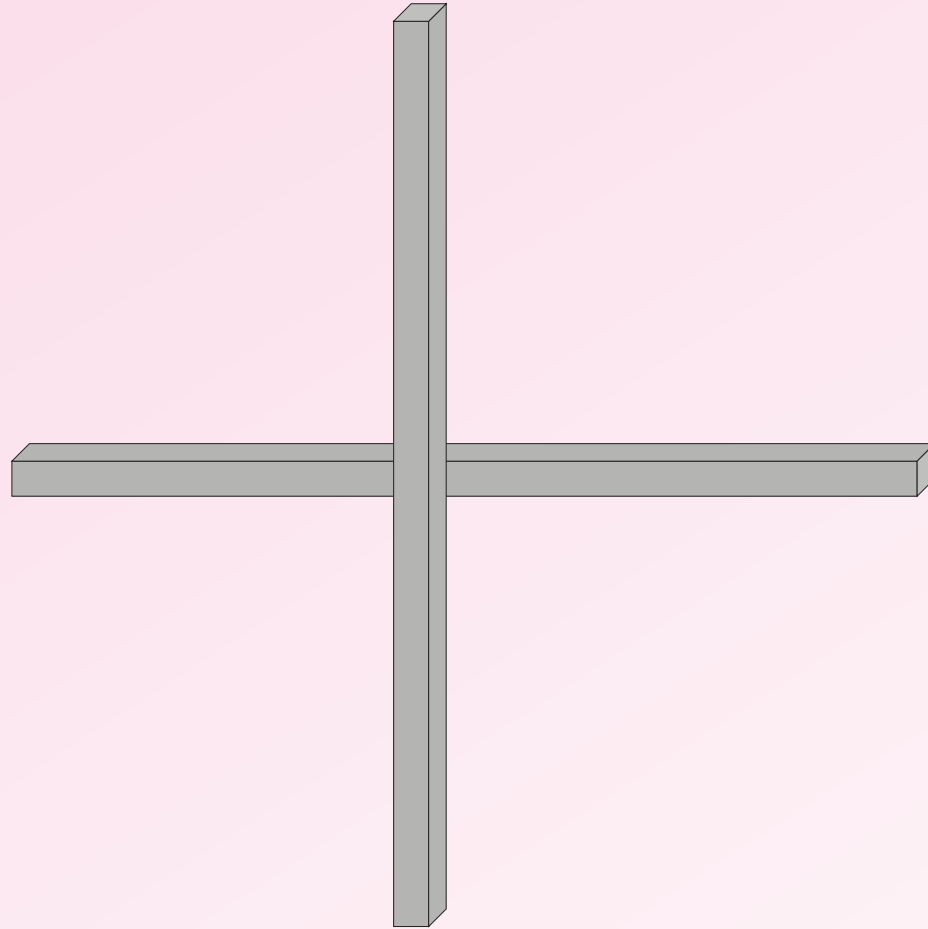
The theorem states that configuration $c(X)$ can be annihilated by a sequence of difference operations.

If $k = 1$ then $c(X)$ is periodic.

More generally, we can prove that $c(X)$ is a sum of k (possibly non-finitary) integral configurations that are periodic.

Corollary. $c(X) = c_1(X) + \dots + c_k(X)$ where $c_i(X)$ is I_i -periodic and integral (but not necessarily finitary).

Example. The 3D counter example



to Nivat's conjecture is a sum of two periodic configurations. It is annihilated by polynomial $(1 - y)(1 - x)$.

Our approach to Nivat's conjecture.

Suppose $P(c, D) \leq |D|$ for some rectangle D .

Then c has annihilating polynomial

$$f(X) = (1 - X^{I_1}) \dots (1 - X^{I_k}).$$

Take the one with smallest k .

If $k = 1$ then c is periodic, so assume that $k \geq 2$.

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Denote $\delta_i(X) = (1 - X^{I_i})$ and $\phi_i(X) = f(X)/\delta_i(X)$.

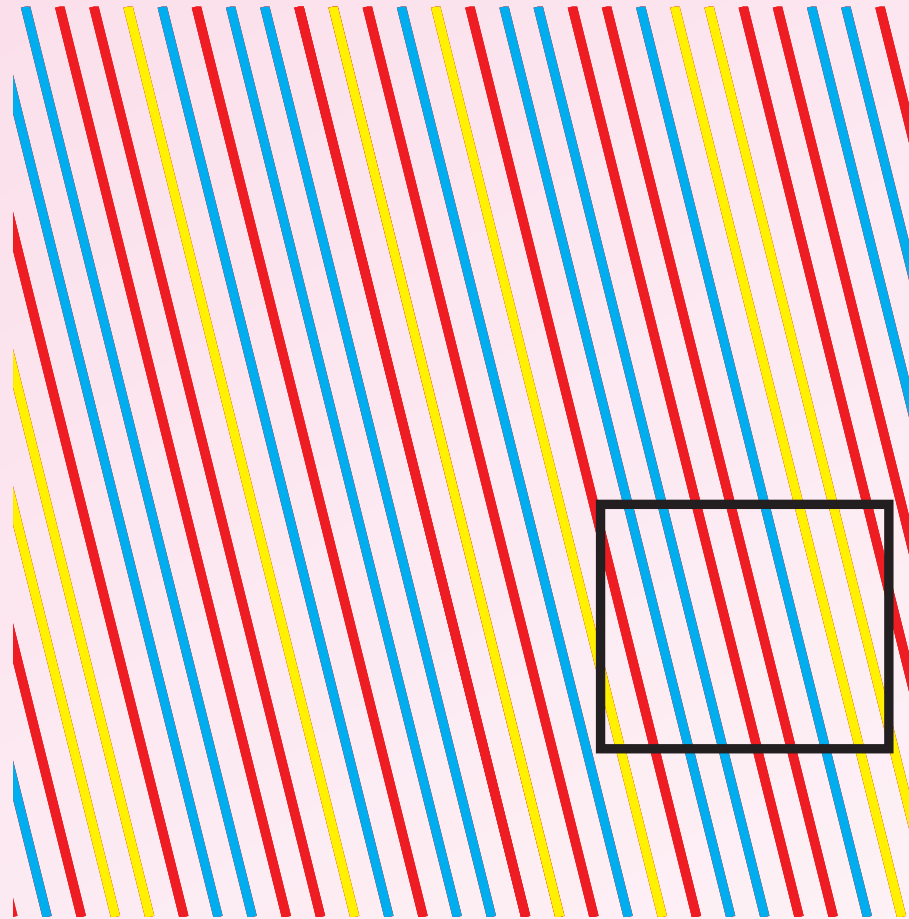
Then $\phi_i(X)c(X)$ is annihilated by $\delta_i(X)$ so it is I_i -periodic. It is not doubly periodic (since otherwise k could be reduced).

Viewing $c(X)$ using filter $\phi_1(X)$:



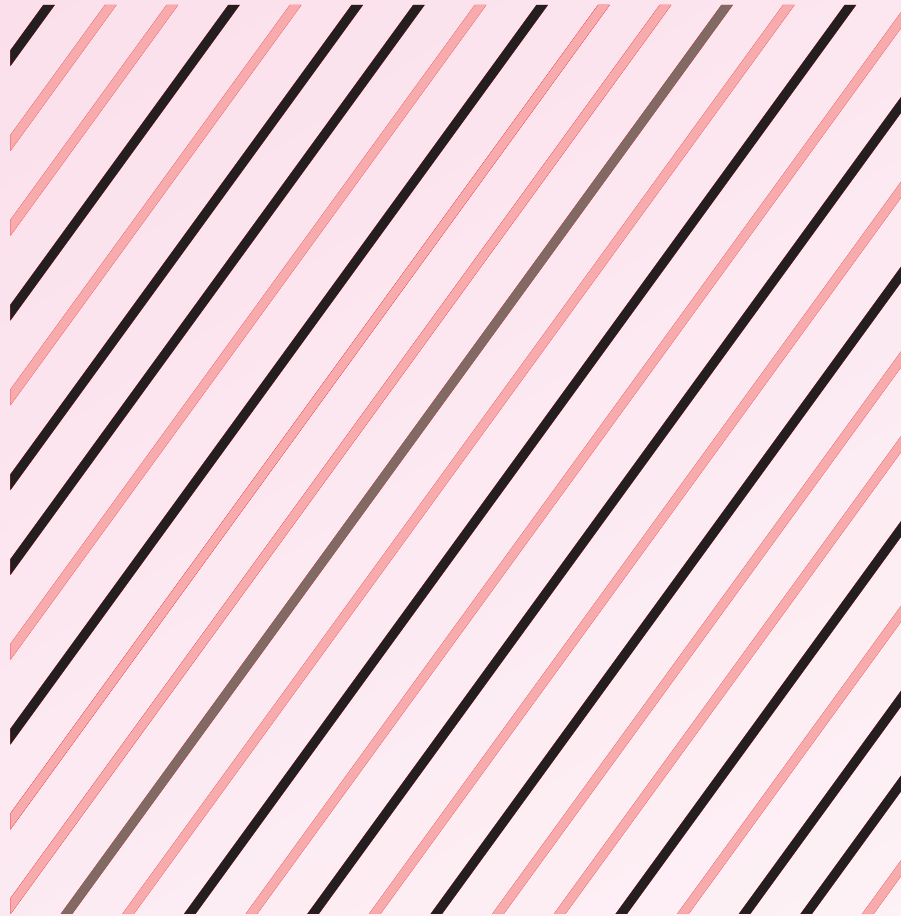
Non-periodic sequence of stripes in the direction I_1 .

Viewing $c(X)$ using filter $\phi_1(X)$:



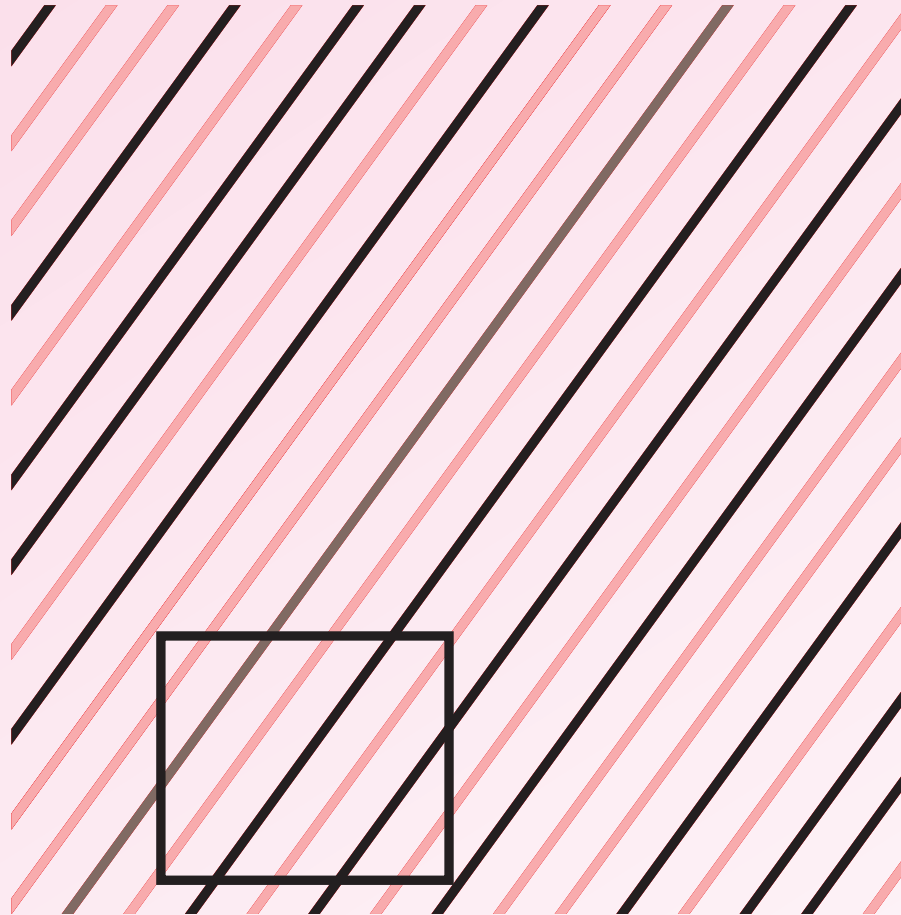
If at least u stripes are visible in every $X \times Y$ rectangle then (by the one-dimensional Morse-Hedlund theorem) there are more than u different $X \times Y$ blocks in $\phi_1(X)c(X)$.

Viewing $c(X)$ using filter $\phi_2(X)$:



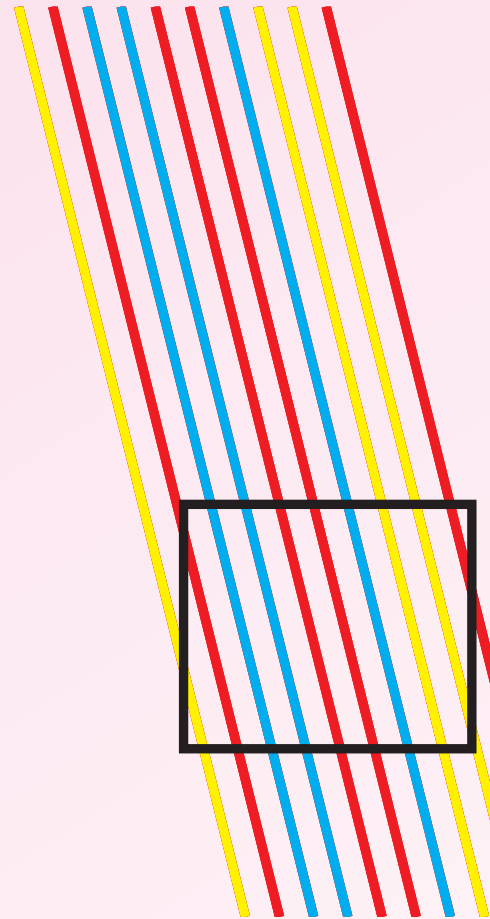
Non-periodic sequence of stripes in a different direction I_2 .

Viewing $c(X)$ using filter $\phi_2(X)$:

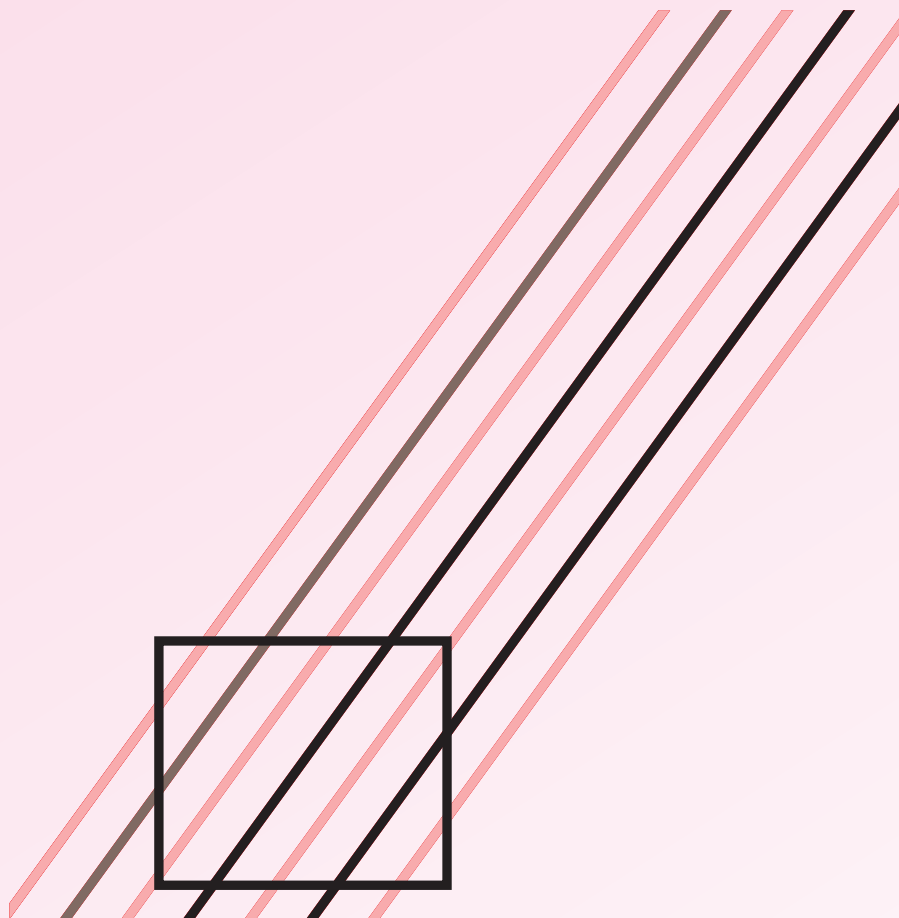


At least v stripes visible in $X \times Y$ rectangles \implies more than v different $X \times Y$ blocks in $\phi_2(X)c(X)$.

Pick any $X \times Y$ pattern from $\phi_1(X)c(X)\dots$



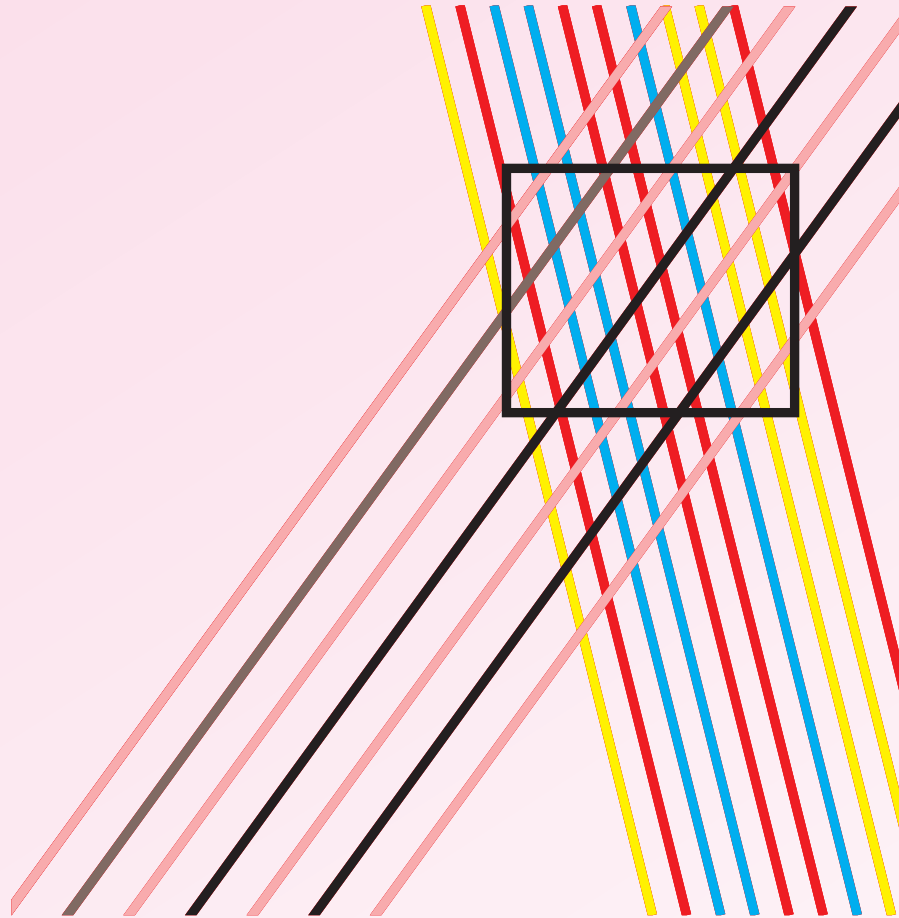
... and any $X \times Y$ pattern from $\phi_2(X)c(X)$.



Directions I_1 and I_2 are different

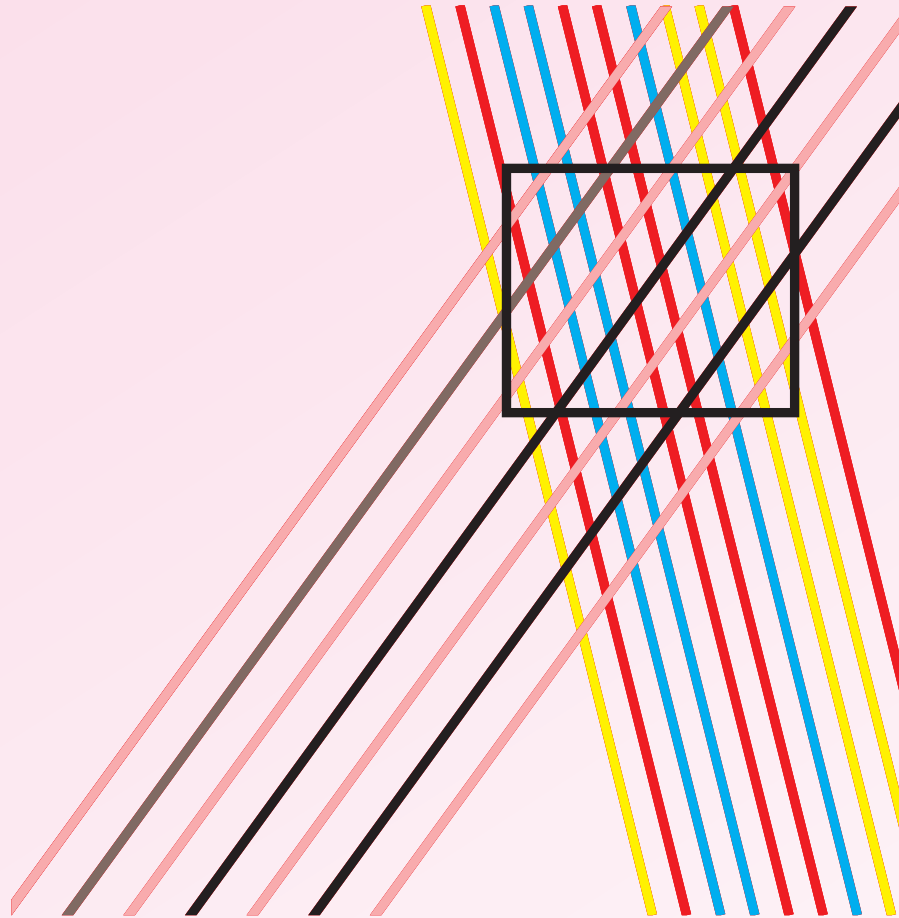


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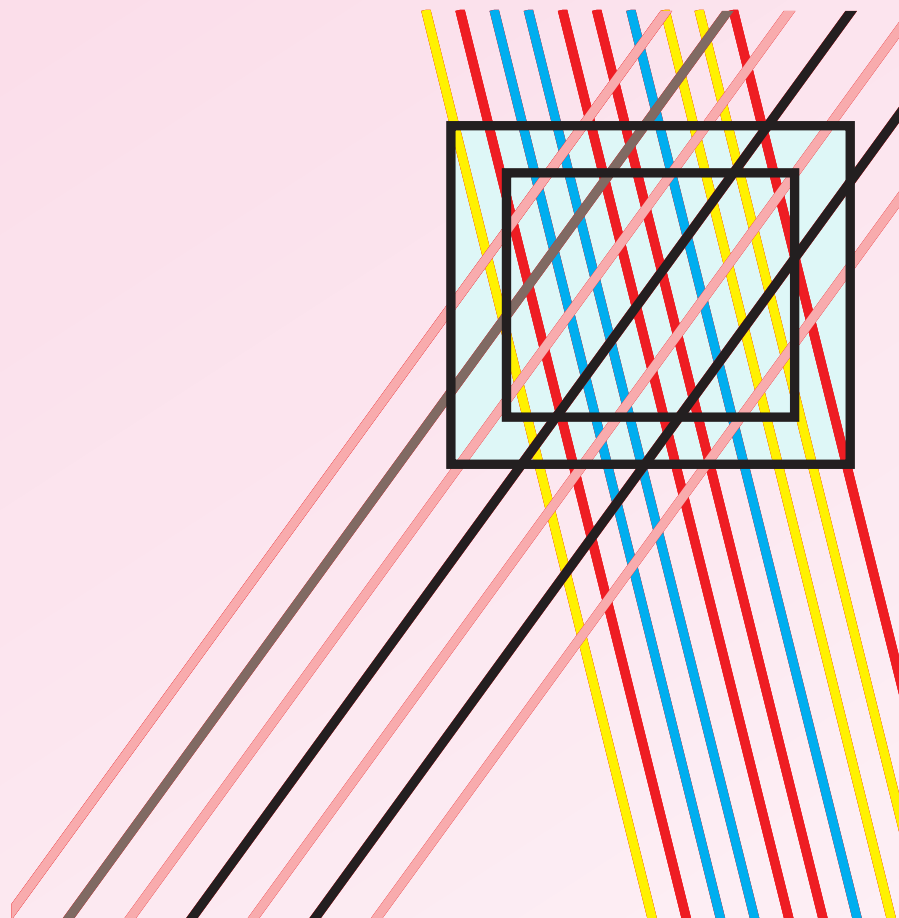
so both patterns can be seen (more or less) in the same position.

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$\implies uv \geq XY$ different pairs of patterns



For some constant r (=radius of filters ϕ_1 and ϕ_2), each $(X + 2r) \times (Y + 2r)$ block of $c(X)$ uniquely determines the corresponding $X \times Y$ blocks in $\phi_1(X)c(X)$ and $\phi_2(X)c(X)$.

$\implies c(X)$ has at least XY patterns of size $(X + 2r) \times (Y + 2r)$.

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Theorem. If c is a non-periodic 2D configuration then $P(c, D) \leq |D|$ can hold only for finitely many rectangles D .

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Thank You